

How an Optimizing Compiler Works

Rewriting code with simple data structures and algorithms

Li Haoyi, Scaladays 12 June 2019

Who Am I

Software Engineer at Databricks

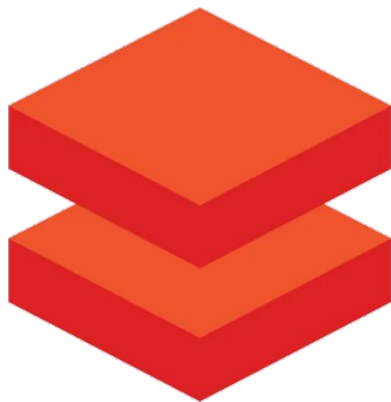
Developer Tools

Lots of Scala internally

Lots of cool technology

Unified Analytics

Hiring in SF and Amsterdam!



Who Am I

Open Source Software Maintainer

com.lihaoyi::sourcecode

com.lihaoyi::fanshi

com.lihaoyi::os-lib

com.lihaoyi::pprint

com.lihaoyi::upack

com.lihaoyi::requests-scala

com.lihaoyi::ammonite

com.lihaoyi::utest

com.lihaoyi::cask

com.lihaoyi::fastparse

com.lihaoyi::ujson

com.lihaoyi::upickle

com.lihaoyi::scalatags

com.lihaoyi::mill



How an Optimizing Compiler Works

Hand Optimizing Some Code

Modelling a Program

Making Inferences and Optimizations

How an Optimizing Compiler Works

Hand Optimizing Some Code

- Type Inference
- Inlining
- Constant Folding
- Dead Code Elimination
- Branch Elimination
- Late Scheduling

Modelling a Program

Making Inferences and Optimizations

Manual Optimizations: Baseline

```
static int main(int n){
    int count = 0, total = 0, multiplied = 0;
    Logger logger = new PrintLogger();
    while(count < n){
        count += 1;
        multiplied *= count;
        if (multiplied < 100) logger.log(count);
        total += ackermann(2, 2);
        total += ackermann(multiplied, n);
        int d1 = ackermann(n, 1);
        total += d1 * multiplied;
        int d2 = ackermann(n, count);
        if (count % 2 == 0) total += d2;
    }
    return total;
}
```

```
// https://en.wikipedia.org/wiki/Ackermann_function
static int ackermann(int m, int n){
    if (m == 0) return n + 1;
    else if (n == 0) return ackermann(m - 1, 1);
    else return ackermann(m - 1, ackermann(m, n - 1));
}

interface Logger{
    public void log(Object a);
}

static class PrintLogger implements Logger{
    public void log(Object a){ System.out.println(a); }
}

static class ErrLogger implements Logger{
    public void log(Object a){ System.err.println(a); }
}
```

Manual Optimizations: Type Inference

```
static int main(int n){
    int count = 0, total = 0, multiplied = 0;
    - Logger logger = new PrintLogger();
    + PrintLogger logger = new PrintLogger();
    while(count < n){
        count += 1;
        multiplied *= count;
        if (multiplied < 100) logger.log(count);
        total += ackermann(2, 2);
        total += ackermann(multiplied, n);
        int d1 = ackermann(n, 1);
        total += d1 * multiplied;
        int d2 = ackermann(n, count);
        if (count % 2 == 0) total += d2;
    }
    return total;
}
```

```
// https://en.wikipedia.org/wiki/Ackermann\_function
static int ackermann(int m, int n){
    if (m == 0) return n + 1;
    else if (n == 0) return ackermann(m - 1, 1);
    else return ackermann(m - 1, ackermann(m, n - 1));
}

interface Logger{
    public void log(Object a);
}

static class PrintLogger implements Logger{
    public void log(Object a){ System.out.println(a); }
}

static class ErrLogger implements Logger{
    public void log(Object a){ System.err.println(a); }
}
```

Manual Optimizations: Inlining

```
static int main(int n){
    int count = 0, total = 0, multiplied = 0;
    PrintLogger logger = new PrintLogger();
    while(count < n){
        count += 1;
        multiplied *= count;
-   if (multiplied < 100) logger.log(count);
+   if (multiplied < 100) System.out.println(count);
        total += ackermann(2, 2);
        total += ackermann(multiplied, n);
        int d1 = ackermann(n, 1);
        total += d1 * multiplied;
        int d2 = ackermann(n, count);
        if (count % 2 == 0) total += d2;
    }
    return total;
}
```

```
// https://en.wikipedia.org/wiki/Ackermann_function
static int ackermann(int m, int n){
    if (m == 0) return n + 1;
    else if (n == 0) return ackermann(m - 1, 1);
    else return ackermann(m - 1, ackermann(m, n - 1));
}

interface Logger{
    public void log(Object a);
}

static class PrintLogger implements Logger{
    public void log(Object a){ System.out.println(a); }
}

static class ErrLogger implements Logger{
    public void log(Object a){ System.err.println(a); }
}
```


Manual Optimizations: Constant Folding

```
static int main(int n){
- int count = 0, total = 0, multiplied = 0;
+ int count = 0, total = 0;
  PrintLogger logger = new PrintLogger();
  while(count < n){
    count += 1;
-   multiplied *= count;
-   if (multiplied < 100) System.out.println(count);
+   if (0 < 100) System.out.println(count);
    total += ackermann(2, 2);
-   total += ackermann(multiplied, n);
+   total += ackermann(0, n);
    int d1 = ackermann(n, 1);
-   total += d1 * multiplied;
    int d2 = ackermann(n, count);
    if (count % 2 == 0) total += d2;
  }
}
```

```
// https://en.wikipedia.org/wiki/Ackermann_function
static int ackermann(int m, int n){
  if (m == 0) return n + 1;
  else if (n == 0) return ackermann(m - 1, 1);
  else return ackermann(m - 1, ackermann(m, n - 1));
}

interface Logger{
  public void log(Object a);
}

static class PrintLogger implements Logger{
  public void log(Object a){ System.out.println(a); }
}

static class ErrLogger implements Logger{
  public void log(Object a){ System.err.println(a); }
}
```

Manual Optimizations: Dead Code Elimination

```
static int main(int n){
    int count = 0, total = 0;
-   PrintLogger logger = new PrintLogger();
    while(count < n){
        count += 1;
        if (0 < 100) System.out.println(count);
        total += ackermann(2, 2);
        total += ackermann(0, n);
-   int d1 = ackermann(n, 1);
        int d2 = ackermann(n, count);
        if (count % 2 == 0) total += d2;
    }
    return total;
}
```

```
// https://en.wikipedia.org/wiki/Ackermann_function
static int ackermann(int m, int n){
    if (m == 0) return n + 1;
    else if (n == 0) return ackermann(m - 1, 1);
    else return ackermann(m - 1, ackermann(m, n - 1));
}
```

```
- interface Logger{
-   public void log(Object a);
- }
- static class PrintLogger implements Logger{
-   public void log(Object a){ System.out.println(a); }
- }
- static class ErrLogger implements Logger{
-   public void log(Object a){ System.err.println(a); }
- }
```

Manual Optimizations: Branch Elimination

```
static int main(int n){
    int count = 0, total = 0;
    while(count < n){
        count += 1;
-   if (0 < 100) System.out.println(count);
+   System.out.println(count);
        total += ackermann(2, 2);
        total += ackermann(0, n);
        int d2 = ackermann(n, count);
        if (count % 2 == 0) total += d2;
    }
    return total;
}
```

```
// https://en.wikipedia.org/wiki/Ackermann\_function
static int ackermann(int m, int n){
    if (m == 0) return n + 1;
    else if (n == 0) return ackermann(m - 1, 1);
    else return ackermann(m - 1, ackermann(m, n - 1));
}
```

Manual Optimizations: Partial Evaluation

```
static int main(int n){
    int count = 0, total = 0;
    while(count < n){
        count += 1;
        System.out.println(count);
        - total += ackermann(2, 2);
        + total += 7;
        - total += ackermann(0, n);
        + total += n + 1;
        int d2 = ackermann(n, count);
        if (count % 2 == 0) total += d2;
    }
    return total;
}
```

```
// https://en.wikipedia.org/wiki/Ackermann\_function
static int ackermann(int m, int n){
    if (m == 0) return n + 1;
    else if (n == 0) return ackermann(m - 1, 1);
    else return ackermann(m - 1, ackermann(m, n - 1));
}
```

Manual Optimizations: Late Scheduling

```
static int main(int n){
    int count = 0, total = 0;
    while(count < n){
        count += 1;
        System.out.println(count);
        total += 7;
        total += n + 1;
-   int d2 = ackermann(n, count);
-   if (count % 2 == 0) total += d2;
+   if (count % 2 == 0) {
+       int d2 = ackermann(n, count);
+       total += d2;
+   }
    }
    return total;
}
```

```
// https://en.wikipedia.org/wiki/Ackermann\_function
static int ackermann(int m, int n){
    if (m == 0) return n + 1;
    else if (n == 0) return ackermann(m - 1, 1);
    else return ackermann(m - 1, ackermann(m, n - 1));
}
```

Manual Optimizations: Final

```
static int main(int n){
    int count = 0, total = 0;
    while(count < n){
        count += 1;
        System.out.println(count);
        total += 7;
        total += n + 1;
        if (count % 2 == 0) {
            int d2 = ackermann(n, count);
            total += d2;
        }
    }
    return total;
}
```

```
// https://en.wikipedia.org/wiki/Ackermann\_function
static int ackermann(int m, int n){
    if (m == 0) return n + 1;
    else if (n == 0) return ackermann(m - 1, 1);
    else return ackermann(m - 1, ackermann(m, n - 1));
}
```

Automated Optimizations

How an Optimizing Compiler Works

Hand Optimizing Some Code

Modelling a Program

- Sourcecode
- Abstract Syntax Trees
- Bytecode
- Dataflow Graphs

Making Inferences and Optimizations

How an Optimizing Compiler Works

Hand Optimizing Some Code

Modelling a Program

- **Sourcecode**
- Abstract Syntax Trees
- Bytecode
- Dataflow Graphs

Making Inferences and Optimizations

Sourcecode

```
""  
static int ackermann(int m, int n){  
    if (m == 0) return n + 1;  
    else if (n == 0) return ackermann(m - 1, 1);  
    else return ackermann(m - 1, ackermann(m, n - 1));  
}  
""
```

Sourcecode

```
""
static int ackermann(int m, int n){
    if (m == 0) return n + 1;
    else if (n == 0) return ackermann(m - 1, 1);
    else return ackermann(m - 1, ackermann(m, n - 1));
}
""
""
static int ackermann(int m, int n){
    // hello I am a comment
    if (m == 0) {
        return n + 1;
    } else if (n == 0) {
        return ackermann(m - 1, 1);
    } else {
        return ackermann(m - 1, ackermann(m, n - 1));
    }
}
""
```

```
""
static int ackermann(int m, int n)      {
    if (m == 0)                          {
        return n + 1;                    }
    else if (n == 0)                      {
        return ackermann(m - 1, 1);      }
    else                                  {
        return ackermann(m - 1, ackermann(m, n - 1));  }}
""
```

How an Optimizing Compiler Works

Hand Optimizing Some Code

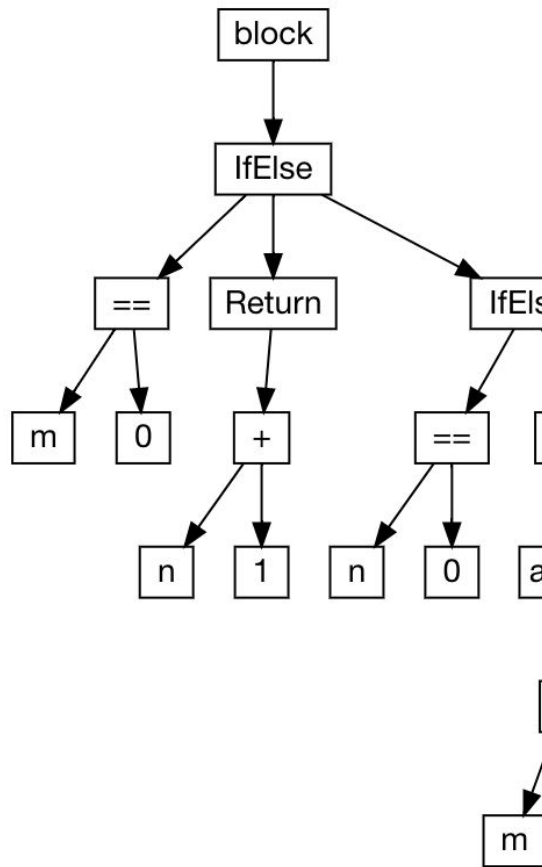
Modelling a Program

- Sourcecode
- **Abstract Syntax Trees**
- Bytecode
- Dataflow Graphs

Making Inferences and Optimizations

Abstract Syntax Trees

```
IfElse(  
  cond = BinOp(Ident("m"), "==", Literal(0)),  
  then = Return(BinOp(Ident("n"), "+", Literal(1))),  
  else = IfElse(  
    cond = BinOp(Ident("n"), "==", Literal(0)),  
    then = Return(Call("ackermann", BinOp(Ident("m"), "-", Literal(1)), Literal(1))),  
    else = Return(  
      Call(  
        "ackermann",  
        BinOp(Ident("m"), "-", Literal(1)),  
        Call("ackermann", Ident("m"), BinOp(Ident("n"), "-", Literal(1)))  
      )  
    )  
  )  
)
```



Abstract Syntax Trees

```
static int ackermannA(int m, int n){  
    int p = n;  
    int q = m;  
    if (q == 0) return p + 1;  
    else if (p == 0) return ackermannA(q - 1, 1);  
    else return ackermannA(q - 1, ackermannA(q, p - 1));  
}
```

```
static int ackermannB(int m, int n){  
    int r = n;  
    int s = m;  
    if (s == 0) return r + 1;  
    else if (r == 0) return ackermannB(s - 1, 1);  
    else return ackermannB(s - 1, ackermannB(s, r - 1));  
}
```

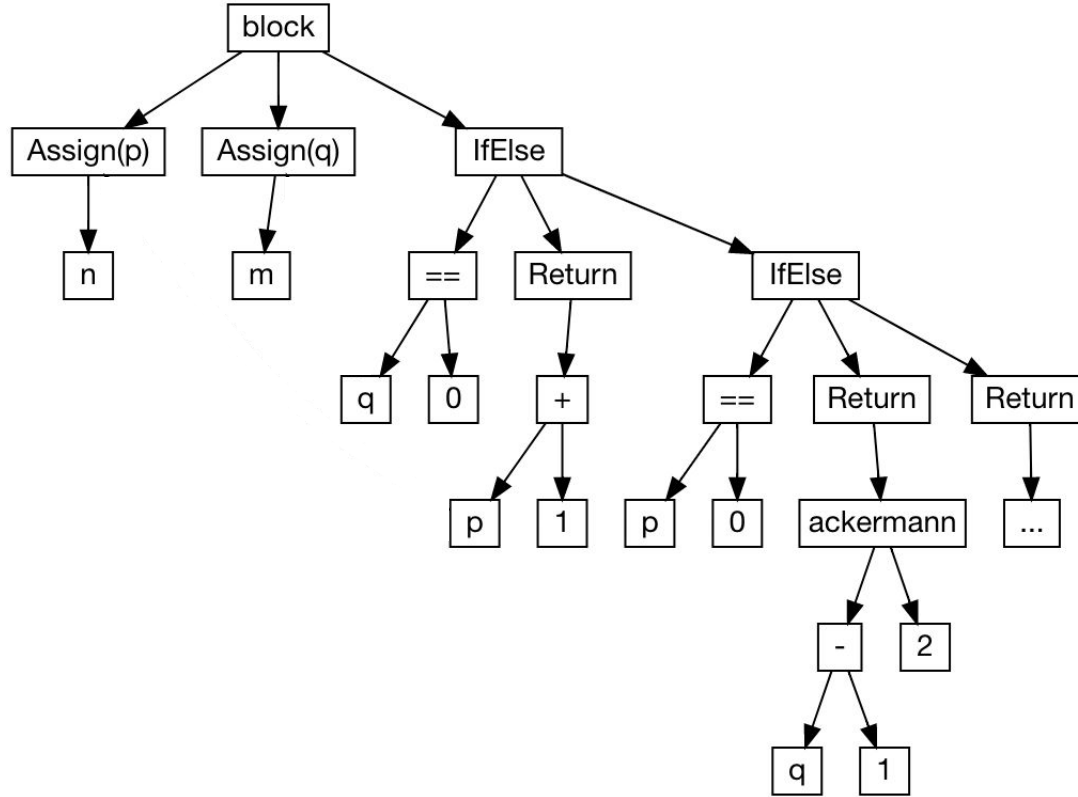
Abstract Syntax Trees

```
Block(  
  Assign("p", Ident("n")),  
  Assign("q", Ident("m")),  
  IfElse(  
    cond = BinOp(Ident("q"), "==", Literal(0)),  
    then = Return(BinOp(Ident("p"), "+", Literal(1))),  
    else = IfElse(  
      cond = BinOp(Ident("p"), "==", Literal(0)),  
      then = Return(Call("ackermann", BinOp(Ident("q"), "-", Literal(1)), Literal(1))),  
      else = Return(  
        Call(  
          "ackermann",  
          BinOp(Ident("q"), "-", Literal(1)),  
          Call("ackermann", Ident("q"), BinOp(Ident("p"), "-", Literal(1)))  
        )  
      )  
    )  
  )  
)
```

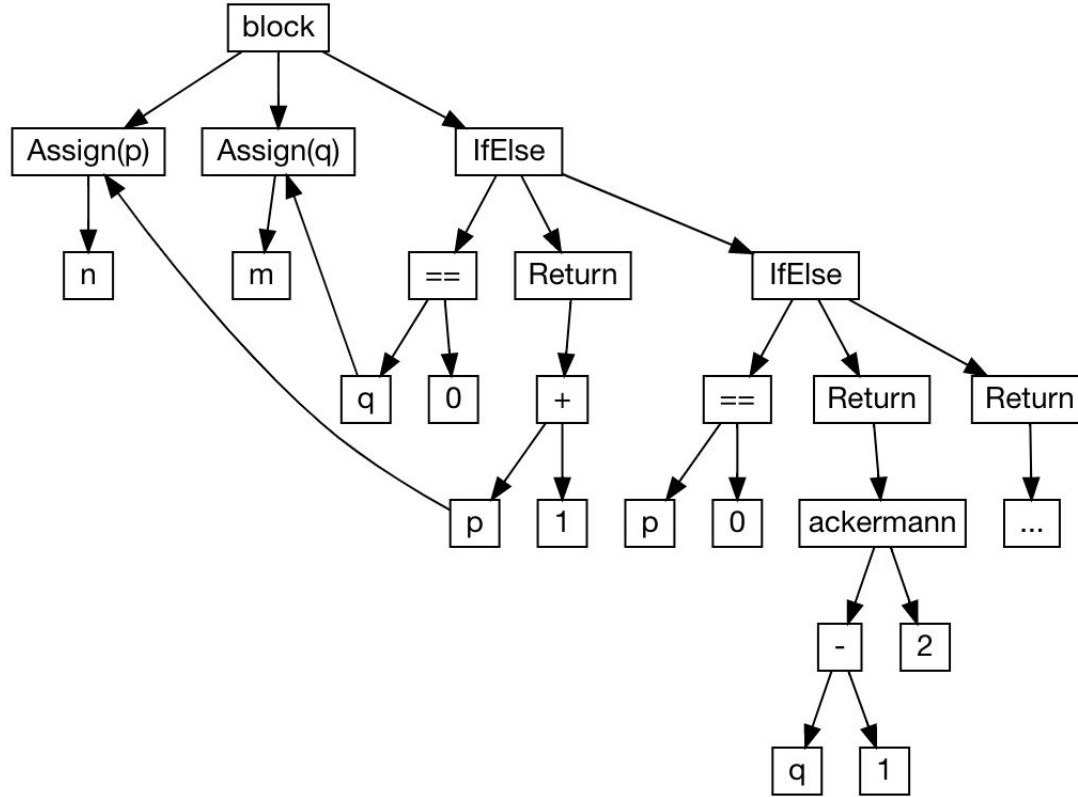
Abstract Syntax Trees

```
Block(  
  Assign("r", Ident("n")),  
  Assign("s", Ident("m")),  
  IfElse(  
    cond = BinOp(Ident("s"), "==", Literal(0)),  
    then = Return(BinOp(Ident("r"), "+", Literal(1))),  
    else = IfElse(  
      cond = BinOp(Ident("r"), "==", Literal(0)),  
      then = Return(Call("ackermann", BinOp(Ident("s"), "-", Literal(1)), Literal(1))),  
      else = Return(  
        Call(  
          "ackermann",  
          BinOp(Ident("s"), "-", Literal(1)),  
          Call("ackermann", Ident("s"), BinOp(Ident("r"), "-", Literal(1)))  
        )  
      )  
    )  
  )  
)
```


Abstract Syntax Trees



Abstract Syntax Trees



How an Optimizing Compiler Works

Hand Optimizing Some Code

Modelling a Program

- Sourcecode
- Abstract Syntax Trees
- **Bytecode**
- Dataflow Graphs

Making Inferences and Optimizations

BYTECODE

```
0: iload_0
1: ifne      8
4: iload_1
5: iconst_1
6: iadd
7: ireturn
8: iload_1
9: ifne     20
12: iload_0
13: iconst_1
14: isub
15: iconst_1
16: invokestatic ackermann:(II)I
19: ireturn
20: iload_0
21: iconst_1
22: isub
23: iload_0
24: iload_1
25: iconst_1
26: isub
27: invokestatic ackermann:(II)I
30: invokestatic ackermann:(II)I
33: ireturn
```

```
static int ackermann(int m, int n){
    if (m == 0) return n + 1;
    else if (n == 0) return ackermann(m - 1, 1);
    else return ackermann(m - 1, ackermann(m, n - 1));
}
```

BYTECODE		LOCALS	STACK	
		a0 a1		
0: iload_0		a0 a1	a0	
1: ifne	8	a0 a1		
4: iload_1		a0 a1	a1	
5: iconst_1		a0 a1	a1 1	static int ackermann(int m, int n){
6: iadd		a0 a1	v1	if (m == 0) return n + 1;
7: ireturn		a0 a1		
8: iload_1		a0 a1	a1	else if (n == 0) return ackermann(m - 1, 1);
9: ifne	20	a0 a1		else return ackermann(m - 1, ackermann(m, n - 1));
12: iload_0		a0 a1	a0	}
13: iconst_1		a0 a1	a0 1	
14: isub		a0 a1	v2	
15: iconst_1		a0 a1	v2 1	
16: invokestatic ackermann:(II)I		a0 a1	v3	
19: ireturn		a0 a1		
20: iload_0		a0 a1	a0	
21: iconst_1		a0 a1	a0 1	
22: isub		a0 a1	v4	
23: iload_0		a0 a1	v4 a0	
24: iload_1		a0 a1	v4 a0 a1	
25: iconst_1		a0 a1	v4 a0 a1 1	
26: isub		a0 a1	v4 a0 v5	
27: invokestatic ackermann:(II)I		a0 a1	v4 v6	
30: invokestatic ackermann:(II)I		a0 a1	v7	
33: ireturn		a0 a1		

BYTECODE

LOCALS

STACK

Bytecode Index	Bytecode	Offset	LOCALS	STACK
0:	iload_0		a0 a1	
1:	ifne	8	a0 a1	a0
4:	iload_1		a0 a1	a1
5:	iconst_1		a0 a1	a1 1
6:	iadd		a0 a1	v1
7:	ireturn		a0 a1	
8:	iload_1		a0 a1	a1
9:	ifne	20	a0 a1	
12:	iload_0		a0 a1	a0
13:	iconst_1		a0 a1	a0 1
14:	isub		a0 a1	v2
15:	iconst_1		a0 a1	v2 1
16:	invokestatic ackermann:(II)I		a0 a1	v3
19:	ireturn		a0 a1	
20:	iload_0		a0 a1	a0
21:	iconst_1		a0 a1	a0 1
22:	isub		a0 a1	v4
23:	iload_0		a0 a1	v4 a0
24:	iload_1		a0 a1	v4 a0 a1
25:	iconst_1		a0 a1	v4 a0 a1 1
26:	isub		a0 a1	v4 a0 v5
27:	invokestatic ackermann:(II)I		a0 a1	v4 v6
30:	invokestatic ackermann:(II)I		a0 a1	v7
30:	invokestatic ackermann2:(I)I		a0 a1	v7
33:	ireturn		a0 a1	

```

static int ackermann(int m, int n){
    if (m == 0) return n + 1;
    else if (n == 0) return ackermann(m - 1, 1);
    else return ackermann(m - 1, ackermann(m, n - 1));
    else return ackermann2(ackermann(m, n - 1));
}

```

BYTECODE

LOCALS

STACK

PC	OpCode	OpArg	LOCALS	STACK
0:	iload_0		a0 a1	
1:	ifne	8	a0 a1	a0
4:	iload_1		a0 a1	a1
5:	iconst_1		a0 a1	a1 1
6:	iadd		a0 a1	v1
7:	ireturn		a0 a1	
8:	iload_1		a0 a1	a1
9:	ifne	20	a0 a1	
12:	iload_0		a0 a1	a0
13:	iconst_1		a0 a1	a0 1
14:	isub		a0 a1	v2
15:	iconst_1		a0 a1	v2 1
16:	invokestatic ackermann:(II)I		a0 a1	v3
19:	ireturn		a0 a1	
20:	iload_0		a0 a1	a0
21:	iconst_1		a0 a1	a0 1
22:	isub		a0 a1	v4
23:	iload_0		a0 a1	v4 a0
24:	iload_1		a0 a1	v4 a0 a1
25:	iconst_1		a0 a1	v4 a0 a1 1
26:	isub		a0 a1	v4 a0 v5
27:	invokestatic ackermann:(II)I		a0 a1	v4 v6
30:	invokestatic ackermann:(II)I		a0 a1	v7
30:	invokestatic ackermann2:(I)I		a0 a1	v7
33:	ireturn		a0 a1	

```

static int ackermann(int m, int n){
    if (m == 0) return n + 1;
    else if (n == 0) return ackermann(m - 1, 1);
    else return ackermann(m - 1, ackermann(m, n - 1));
    else return ackermann2(ackermann(m, n - 1));
}

```


How an Optimizing Compiler Works

Hand Optimizing Some Code

Modelling a Program

- Sourcecode
- Abstract Syntax Trees
- Bytecode
- **Dataflow Graphs**

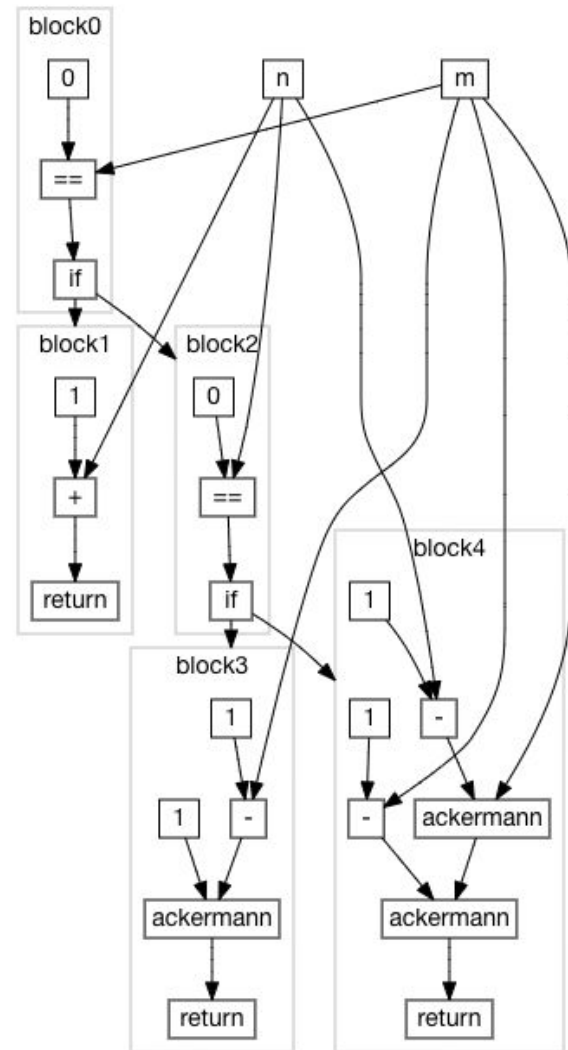
Making Inferences and Optimizations

Dataflow Graphs

```
static int ackermann(int m, int n){
    if (m == 0) return n + 1;
    else if (n == 0) return ackermann(m - 1, 1);
    else return ackermann(m - 1, ackermann(m, n - 1));
}

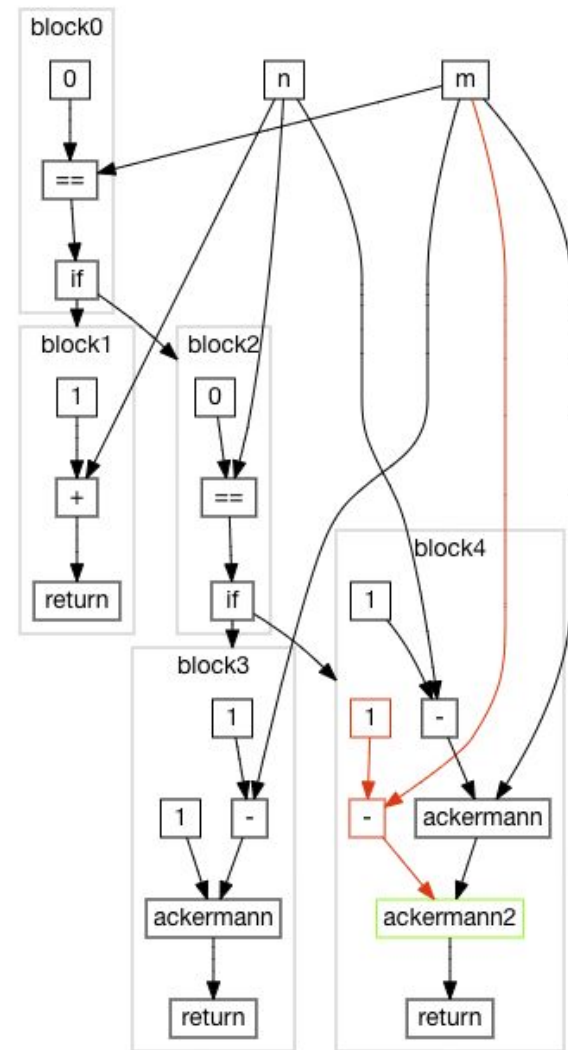
static int ackermannA(int m, int n){
    int p = n;
    int q = m;
    if (q == 0) return p + 1;
    else if (p == 0) return ackermannA(q - 1, 1);
    else return ackermannA(q - 1, ackermannA(q, p - 1));
}

static int ackermannB(int m, int n){
    int r = n;
    int s = m;
    if (s == 0) return r + 1;
    else if (r == 0) return ackermannB(s - 1, 1);
```



Dataflow Graphs

```
static int ackermann(int m, int n){  
    if (m == 0) return n + 1;  
    else if (n == 0) return ackermann(m - 1, 1);  
    else return ackermann(m - 1, ackermann(m, n - 1));  
    else return ackermann2(ackermann(m, n - 1));  
}
```



How an Optimizing Compiler Works

Hand Optimizing Some Code

Modelling a Program

Inferences and Optimizations

- Type Inference & Constant Folding
- Inter-Procedural Inference
- Recursive Inter-Procedural Inference
- Liveness & Reachability Analysis

How an Optimizing Compiler Works

Hand Optimizing Some Code

Modelling a Program

Inferences and Optimizations

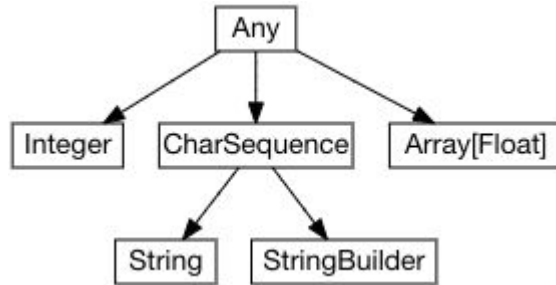
- **Type Inference & Constant Folding**
- Inter-Procedural Inference
- Recursive Inter-Procedural Inference
- Liveness & Reachability Analysis

Type Inference & Constant Folding

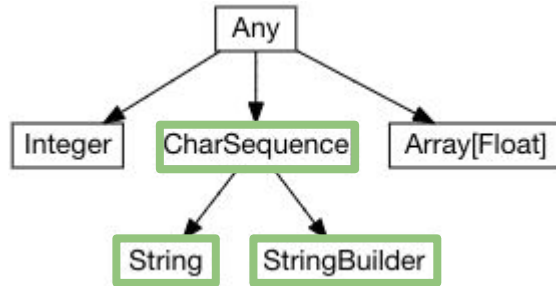
What do we know about a value?

- Is it an Integer? String? Array[Float]? PrintLogger?
- Is it a CharSequence, which could be either a String or a StringBuilder?
- Is it Any, meaning we don't know anything about it?

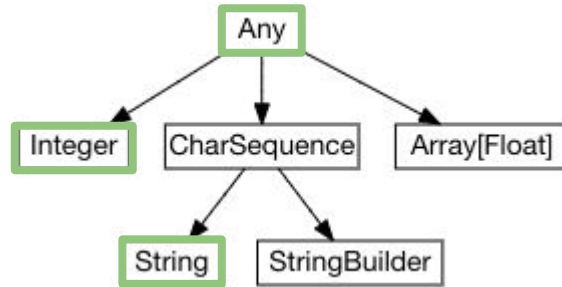
Type Lattices



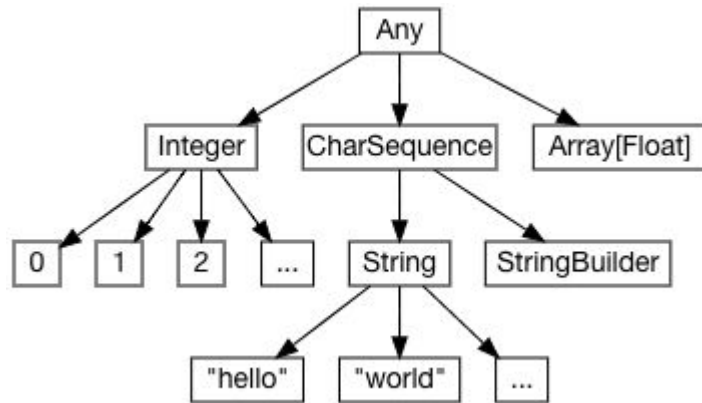
Type Lattices



Type Lattices

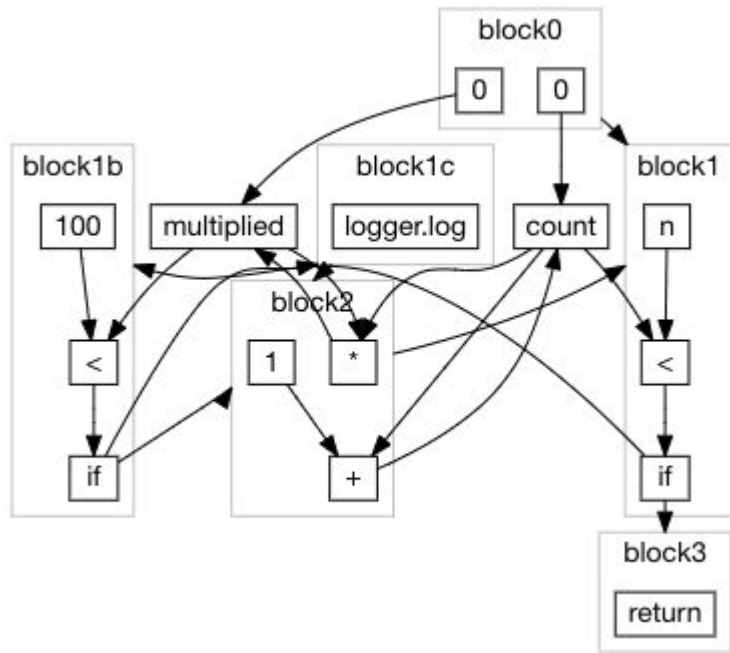


Type Lattices



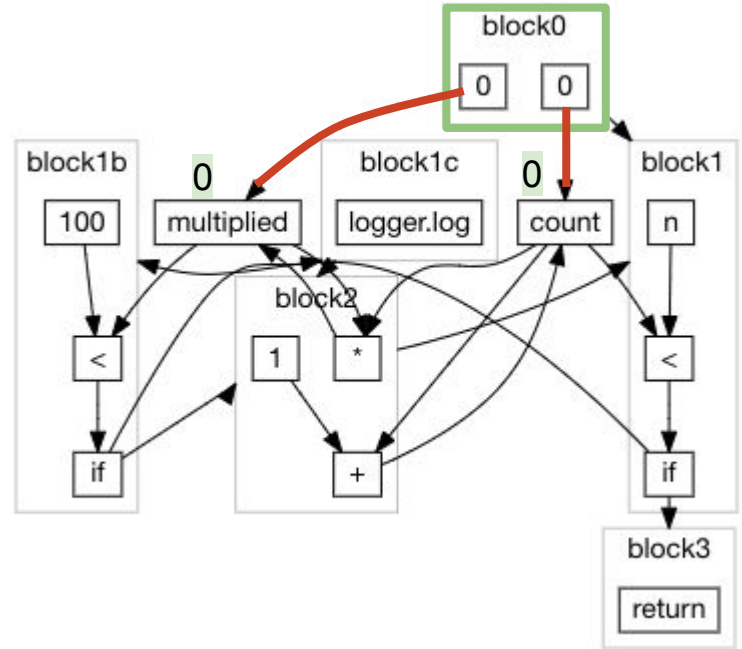
Inferring Values on the Dataflow Graph

```
static int main(int n){  
    int count = 0, multiplied = 0;  
    while(count < n){  
        if (multiplied < 100) logger.log(count);  
        count += 1;  
        multiplied *= count;  
    }  
    return ...;  
}
```



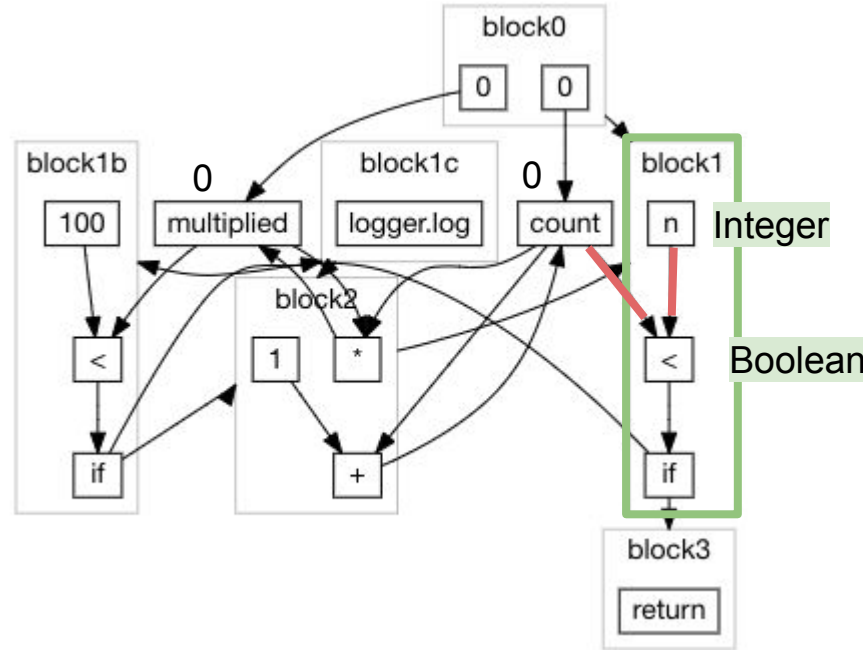
Inferring Values on the Dataflow Graph

```
static int main(int n){  
    int count = 0, multiplied = 0;  
    while(count < n){  
        if (multiplied < 100) logger.log(count);  
        count += 1;  
        multiplied *= count;  
    }  
    return ...;  
}
```



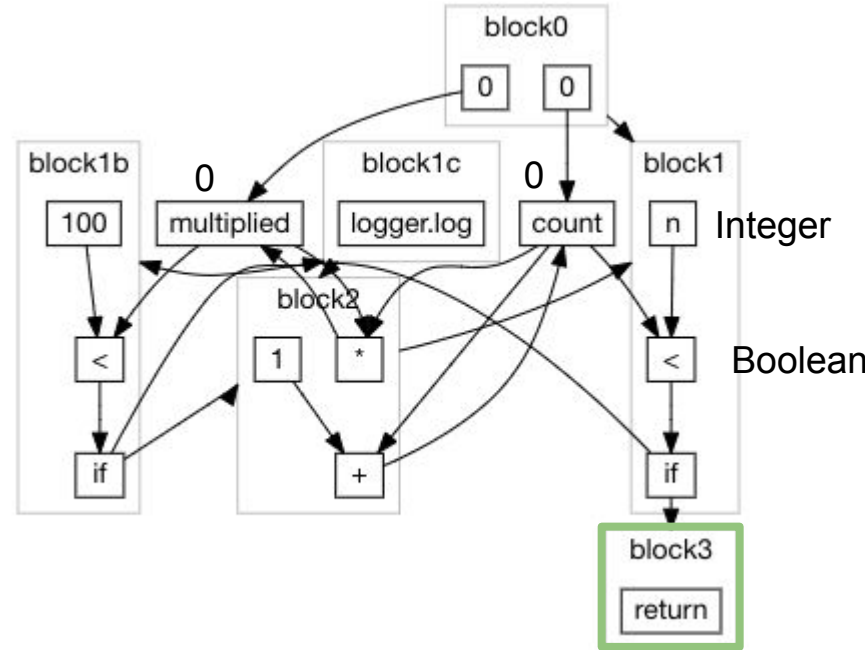
Inferring Values on the Dataflow Graph

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    while(count < n){  
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        count += 1;  
        multiplied *= count;  
    }  
    return ...;  
}
```



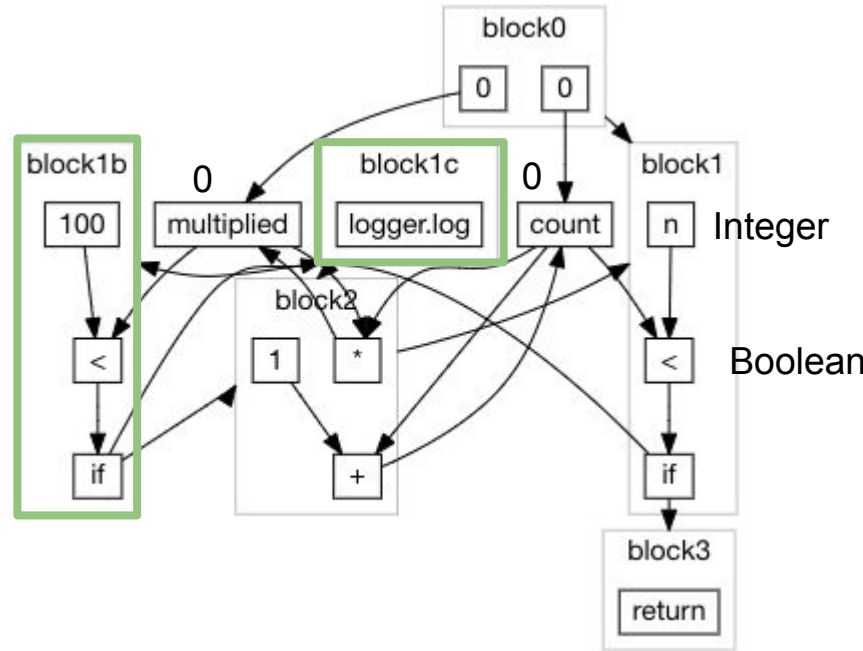
Inferring Values on the Dataflow Graph

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        count += 1;  
        multiplied *= count;  
    }  
    return ...;  
}
```



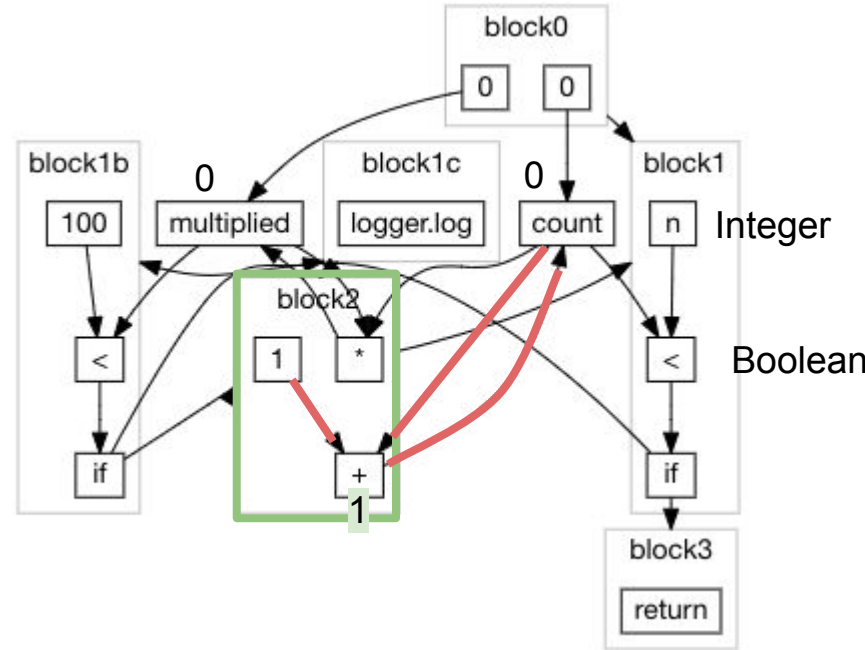
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        count += 1;  
        multiplied *= count;  
    }  
    return ...;  
}
```



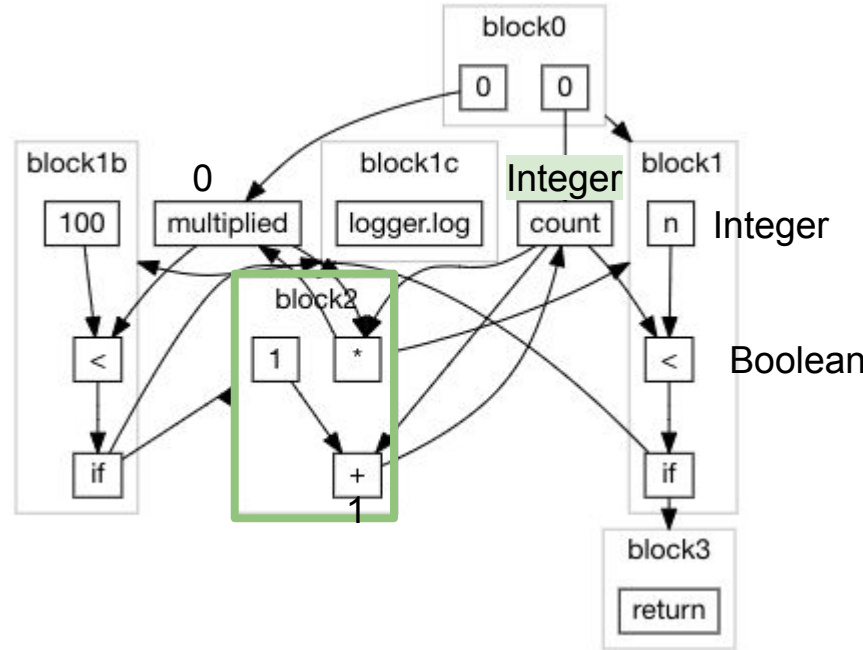
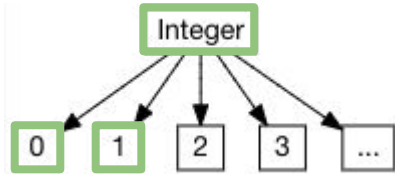
Inferring Values on the Dataflow Graph

```
static int main(int n){  
    int count = 0, multiplied = 0;  
    while(count < n){  
        if (multiplied < 100) logger.log(count);  
        count += 1;  
        multiplied *= count;  
    }  
    return ...;  
}
```



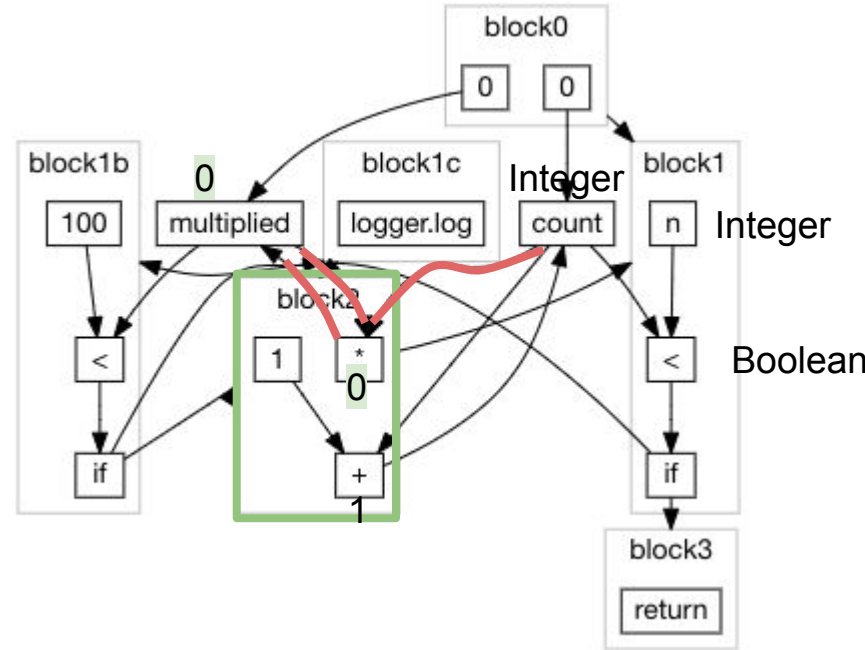
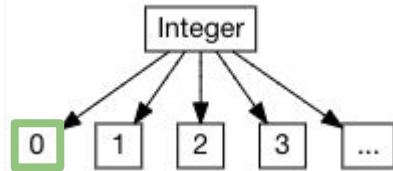
Inferring Values on the Dataflow Graph

```
static int main(int n){  
    int count = 0, multiplied = 0;  
    while(count < n){  
        if (multiplied < 100) logger.log(count);  
        count += 1;  
        multiplied *= count;  
    }  
    return ...;  
}
```



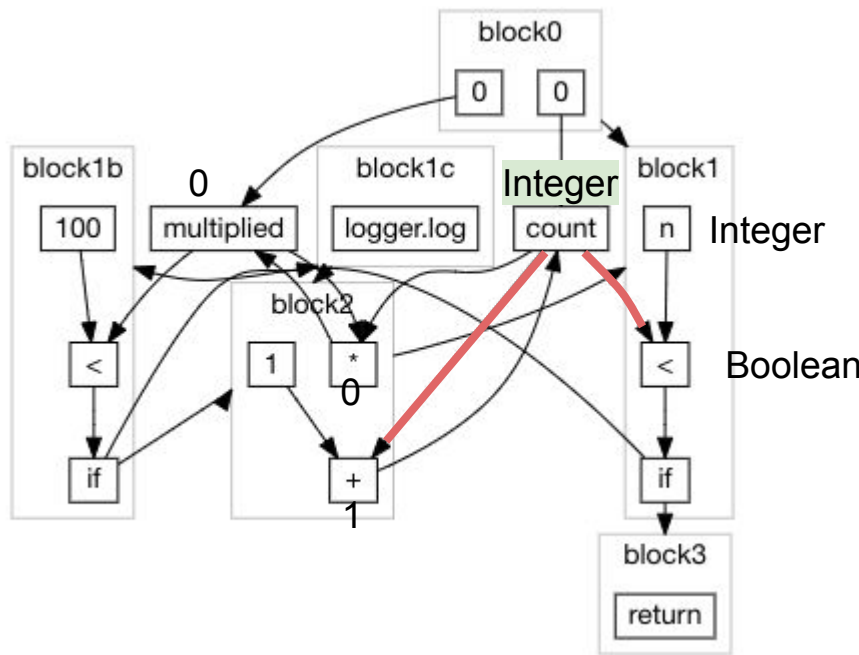
Inferring Values on the Dataflow Graph

```
static int main(int n){  
    int count = 0, multiplied = 0;  
    while(count < n){  
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        count += 1;  
        multiplied *= count;  
    }  
    return ...;  
}
```



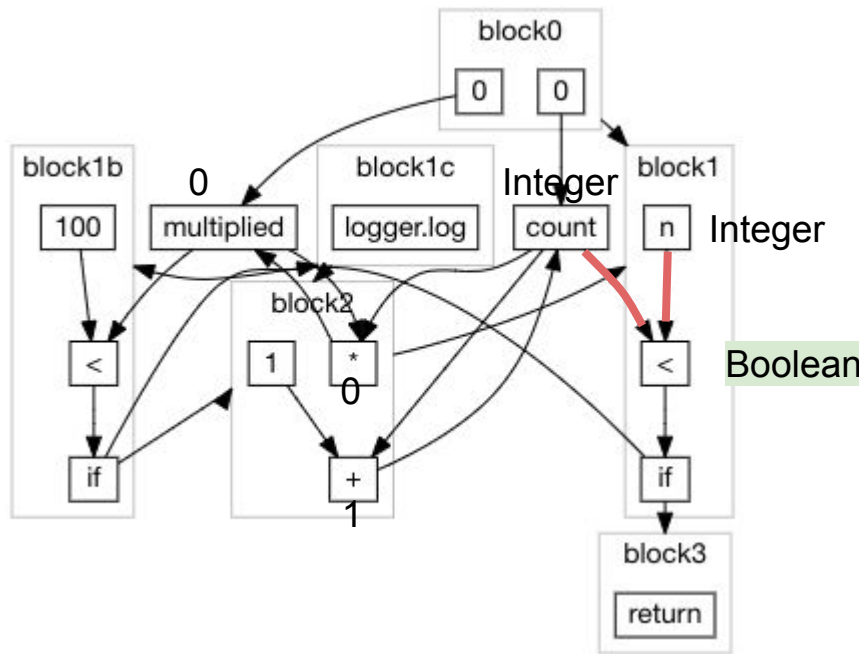
Inferring Values on the Dataflow Graph

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static int main(int n){  
    int count = 0, multiplied = 0;  
    while(count < n){  
        if (multiplied < 100) logger.log(count);  
        count += 1;  
        multiplied *= count;  
    }  
    return ...;  
}
```



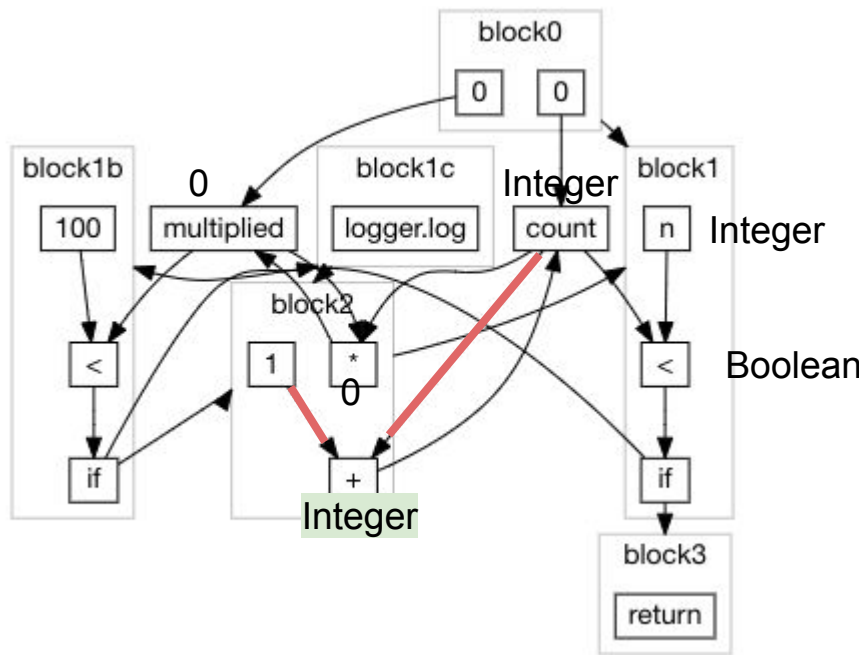
Inferring Values on the Dataflow Graph

```
static int main(int n){  
    int count = 0, multiplied = 0;  
    while(count < n){  
        if (multiplied < 100) logger.log(count);  
        count += 1;  
        multiplied *= count;  
    }  
    return ...;  
}
```



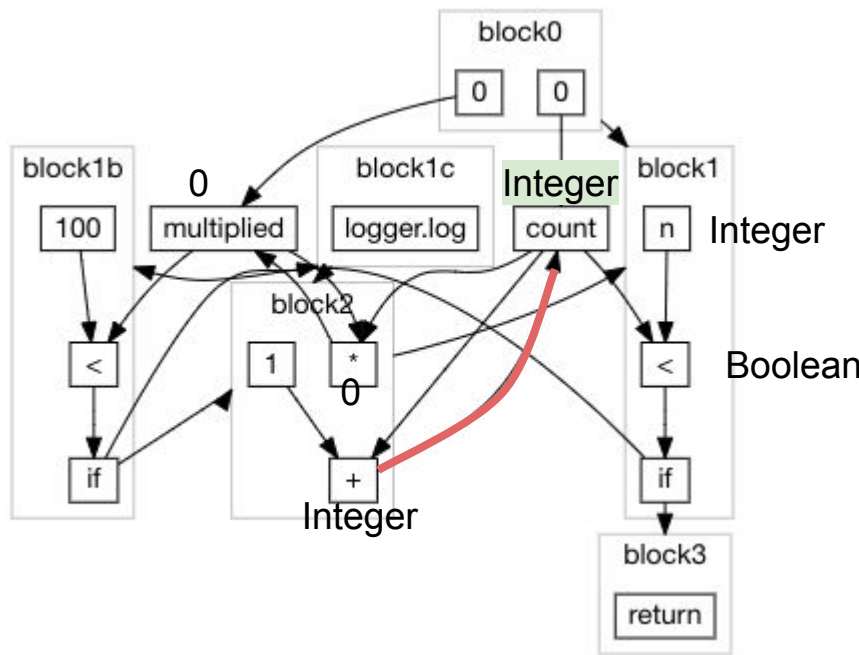
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        count += 1;  
        multiplied *= count;  
    }  
    return ...;  
}
```



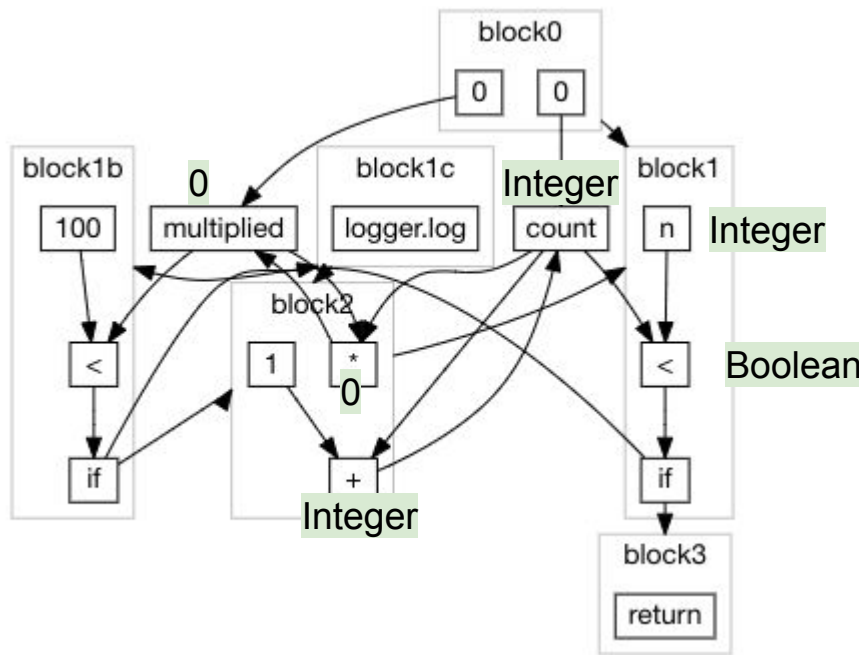
Inferring Values on the Dataflow Graph

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        count += 1;
        multiplied *= count;
    }
    return ...;
}
```



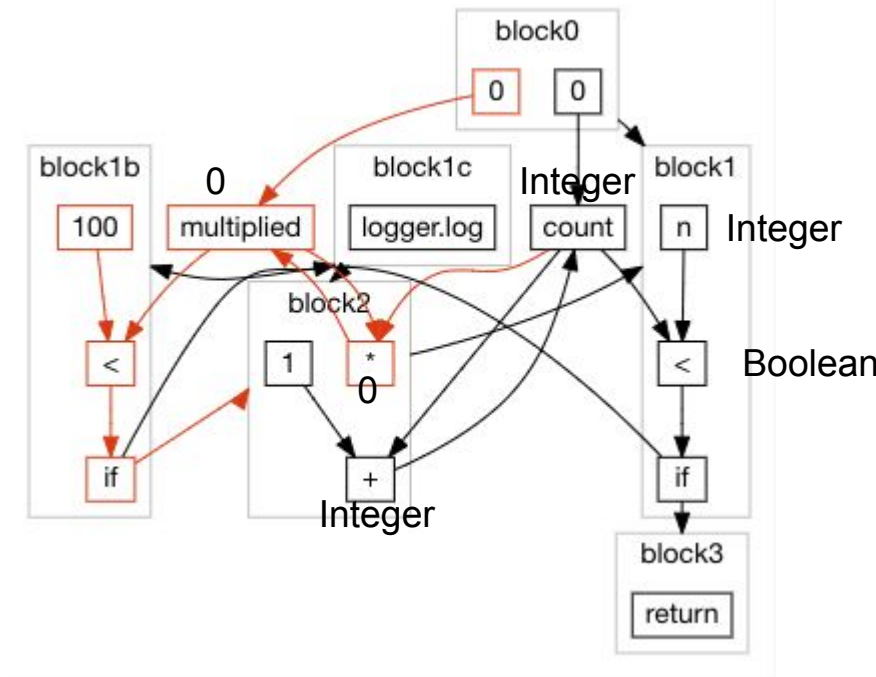
Inferring Values on the Dataflow Graph

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static int main(int n){  
    int count = 0, multiplied = 0;  
    while(count < n){  
        if (multiplied < 100) logger.log(count);  
        count += 1;  
        multiplied *= count;  
    }  
    return ...;  
}
```



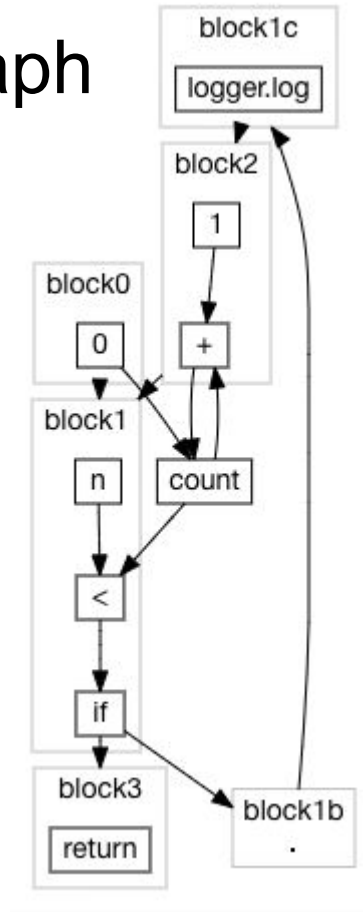
Inferring Values on the Dataflow Graph

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    while(count < n){  
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        count += 1;  
        multiplied *= count;  
    }  
    return ...;  
}
```



Inferring Values on the Dataflow Graph

```
static int main(int n){  
    int count = 0;  
    while(count < n){  
        logger.log(count);  
        count += 1;  
    }  
    return ...;  
}
```



How an Optimizing Compiler Works

Hand Optimizing Some Code

Modelling a Program

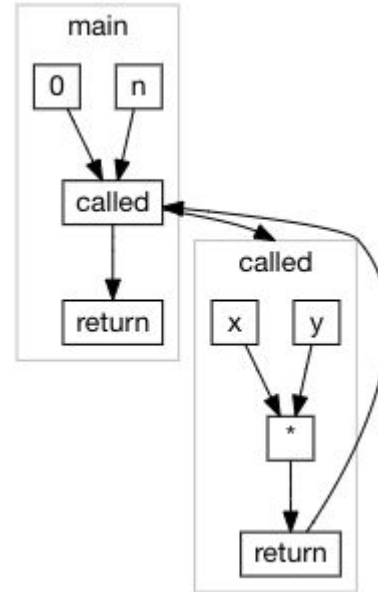
Inferences and Optimizations

- Type Inference & Constant Folding
- **Inter-Procedural Inference**
- Recursive Inter-Procedural Inference
- Liveness & Reachability Analysis

Inter-Procedural Inference

```
static int main(int n){  
    return called(0, n);  
}
```

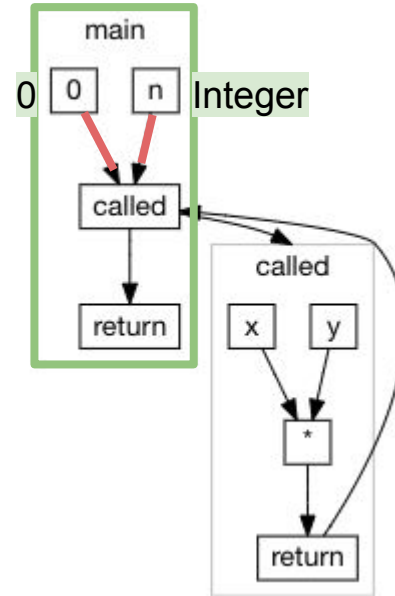
```
static int called(int x, int y){  
    return x * y;  
}
```



Inter-Procedural Inference

```
static int main(int n){  
    return called(0, n);  
}
```

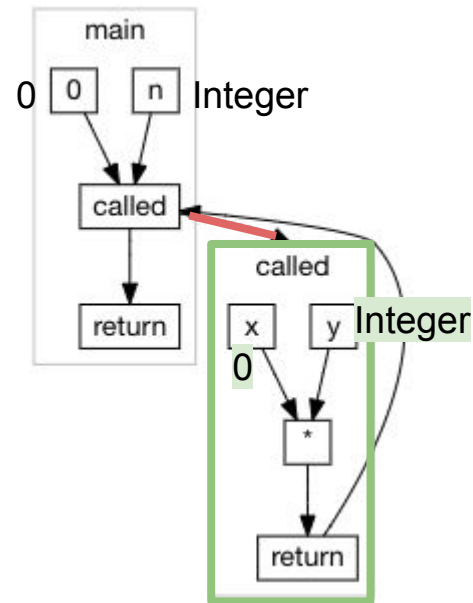
```
static int called(int x, int y){  
    return x * y;  
}
```



Inter-Procedural Inference

```
static int main(int n){  
    return called(0, n);  
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```

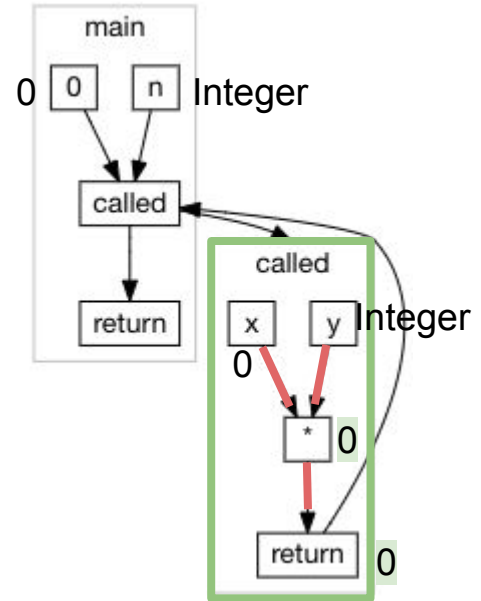
```
static int called(int x, int y){  
    return x * y;  
}
```



Inter-Procedural Inference

```
static int main(int n){  
    return called(0, n);  
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```

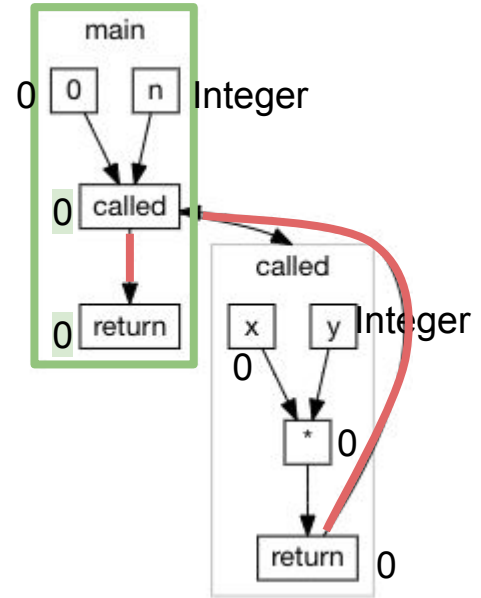
```
static int called(int x, int y){  
    return x * y;  
}
```



Inter-Procedural Inference

```
static int main(int n){  
    return called(0, n);  
}
```

```
static int called(int x, int y){  
    return x * y;  
}
```

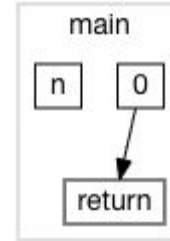


Inter-Procedural Inference

```
static int main(int n){  
    return called(n, 0);  
}
```

```
static int called(int x, int y){  
    return x * y;  
}
```

```
static int main(int n){  
    return 0;  
}
```



How an Optimizing Compiler Works

Hand Optimizing Some Code

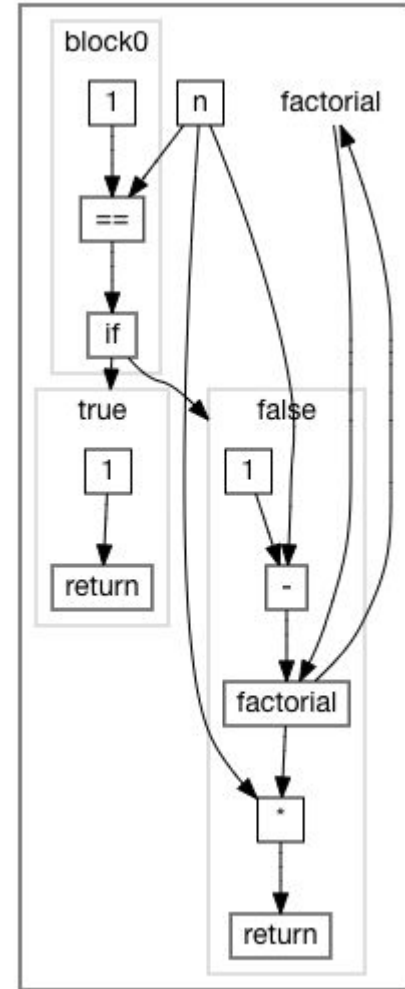
Modelling a Program

Inferences and Optimizations

- Type Inference & Constant Folding
- Inter-Procedural Inference
- **Recursive Inter-Procedural Inference**
- Liveness & Reachability Analysis

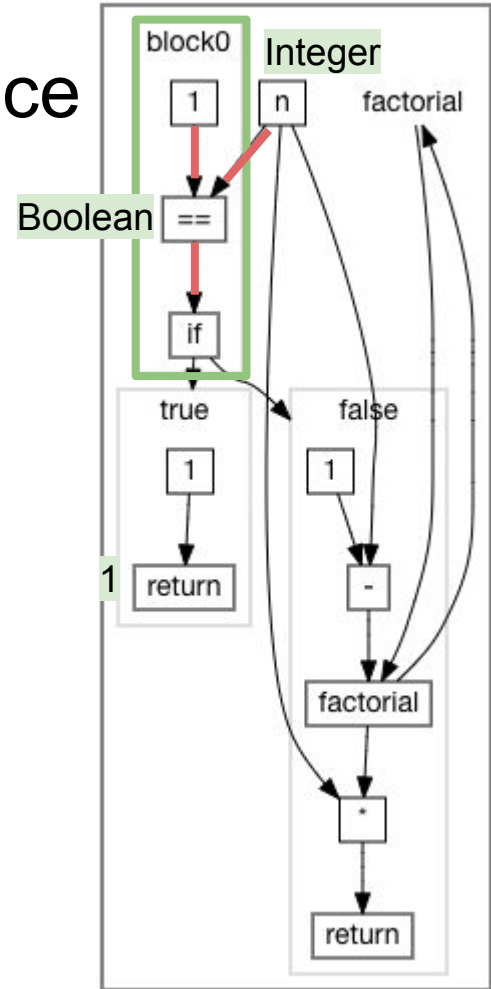
Recursive Inter-Procedural Inference

```
public static Any factorial(int n) {  
    if (n == 1) {  
        return 1;  
    } else {  
        return n * factorial(n - 1);  
    }  
}
```



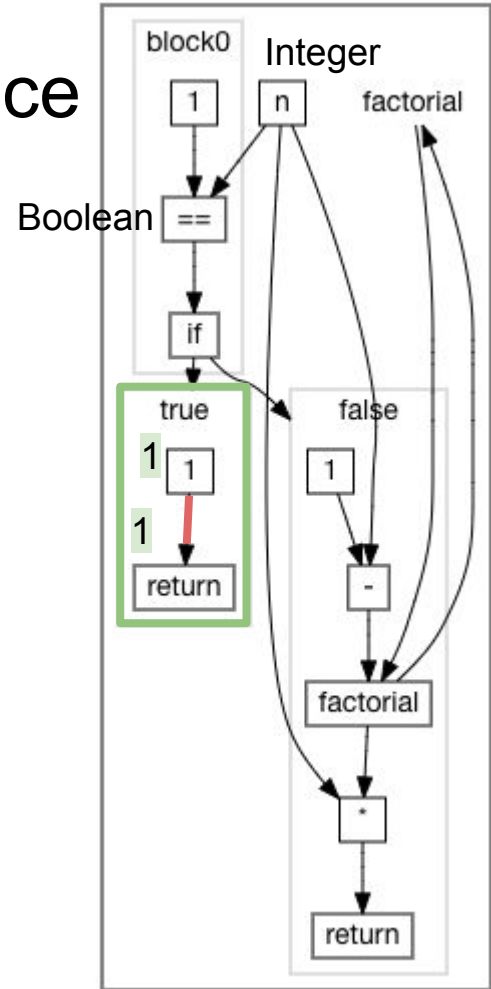
Recursive Inter-Procedural Inference

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    } else {  
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}
```



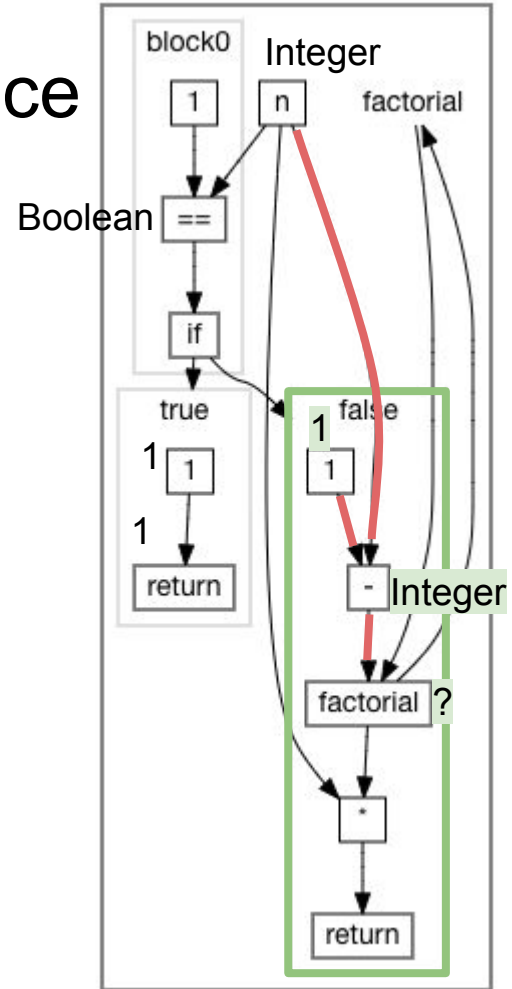
Recursive Inter-Procedural Inference

```
public static Any factorial(int n) {  
    if (n == 1) {  
        return 1;  
    } else {  
        return n * factorial(n - 1);  
    }  
}
```



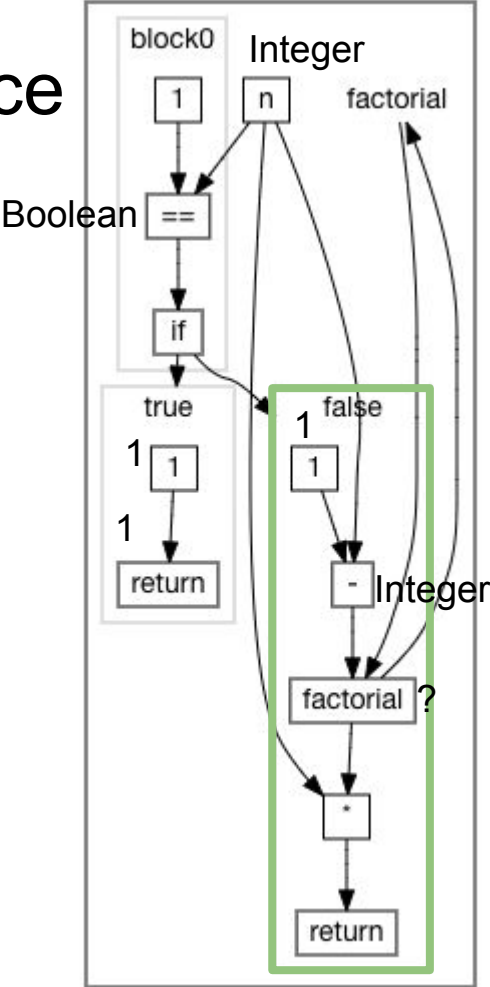
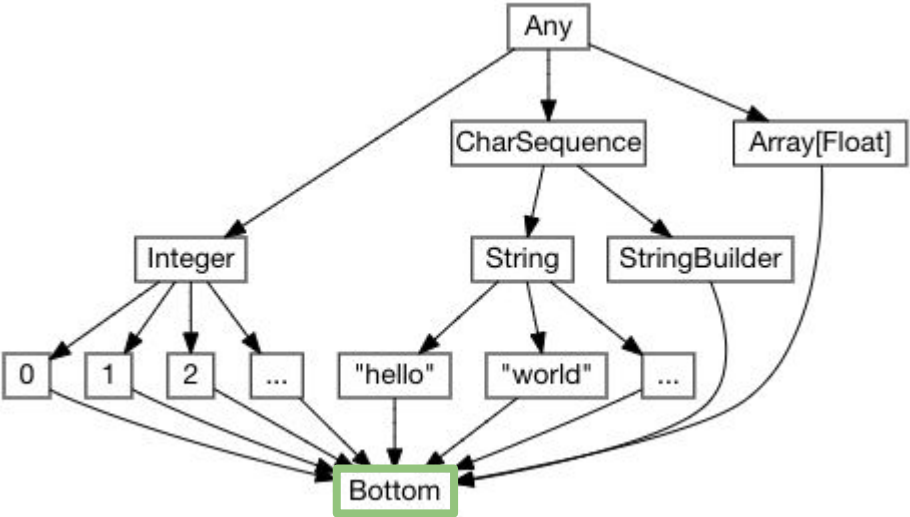
Recursive Inter-Procedural Inference

```
public static Any factorial(int n) {  
    if (n == 1) {  
        return 1;  
    } else {  
        return n * factorial(n - 1);  
    }  
}
```



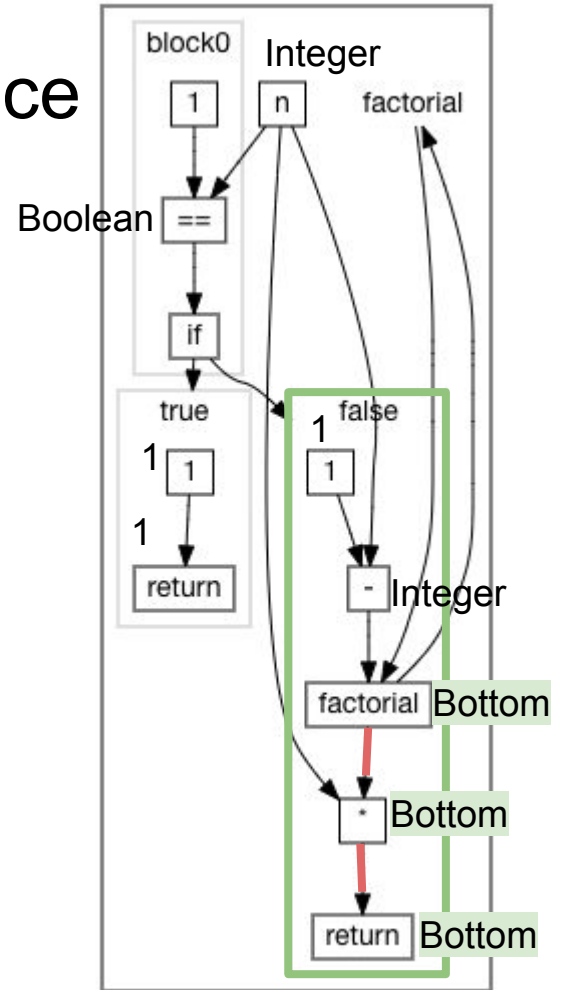
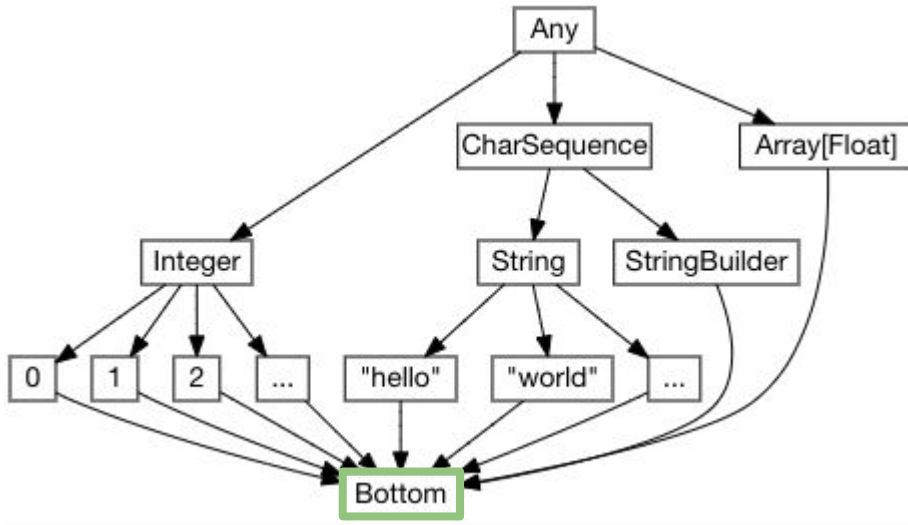
Recursive Inter-Procedural Inference

```
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    if (n == 1) {  
        return 1;  
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        return n * factorial(n - 1);  
    }  
}
```



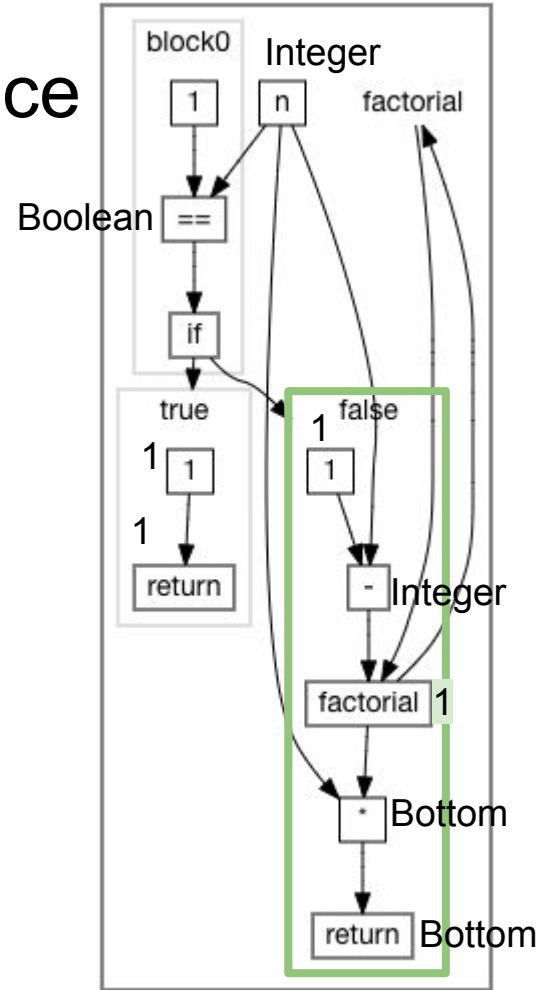
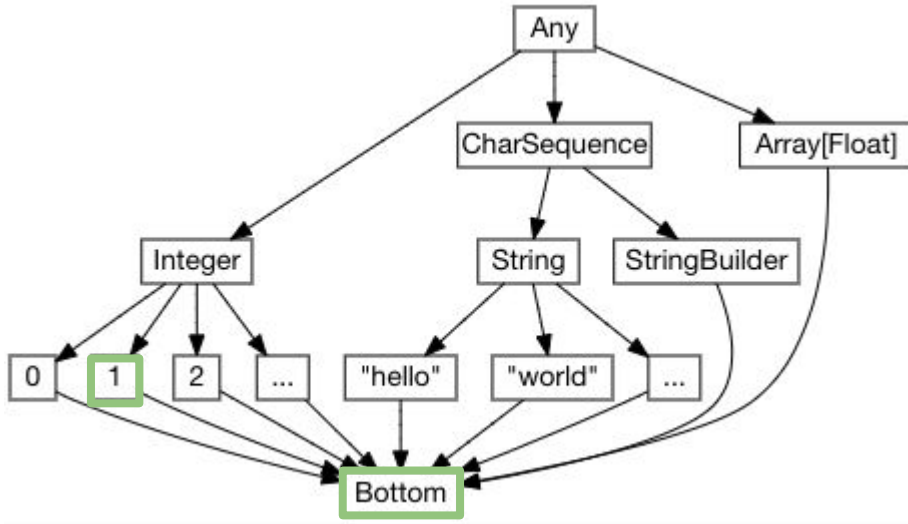
Recursive Inter-Procedural Inference

```
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    if (n == 1) {  
        return 1;  
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    }  
}
```



Recursive Inter-Procedural Inference

```
public static Any factorial(int n) {  
    if (n == 1) {  
        return 1;  
    } else {  
        return n * factorial(n - 1);  
    }  
}
```

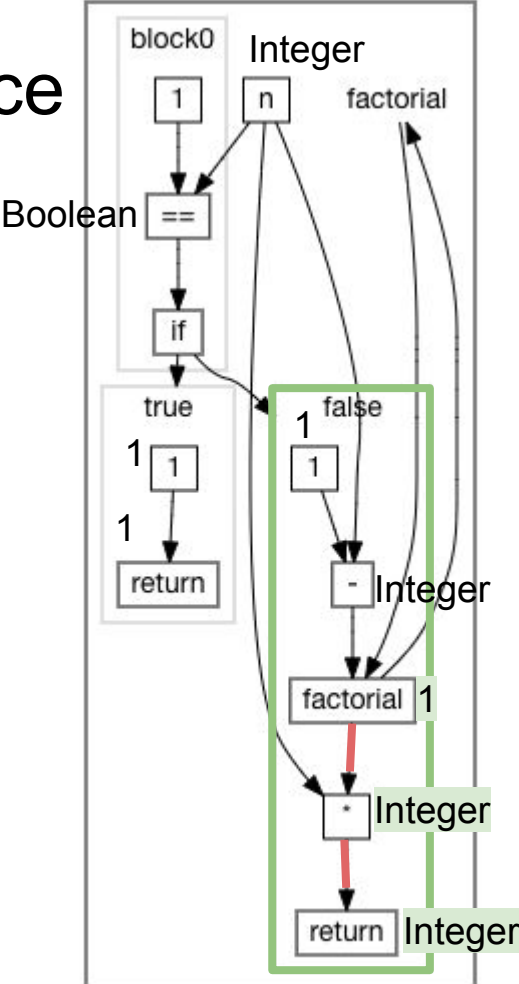
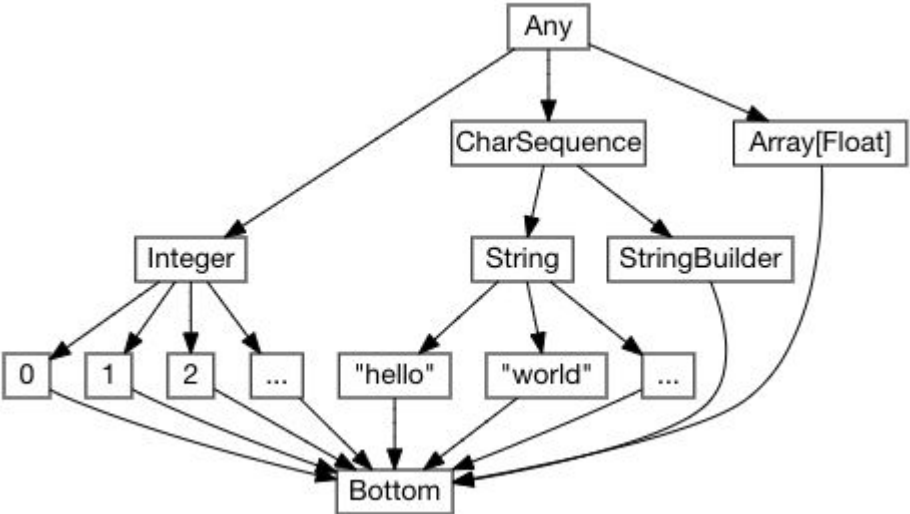


Recursive Inter-Procedural Inference

```

public static Any factorial(int n) {
    if (n == 1) {
        return 1;
    } else {
        return n * factorial(n - 1);
    }
}

```

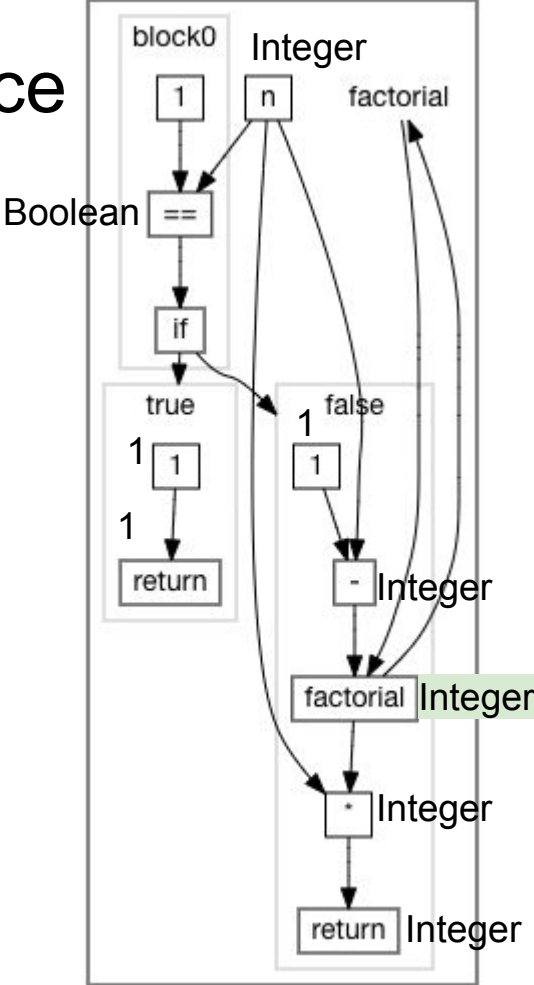
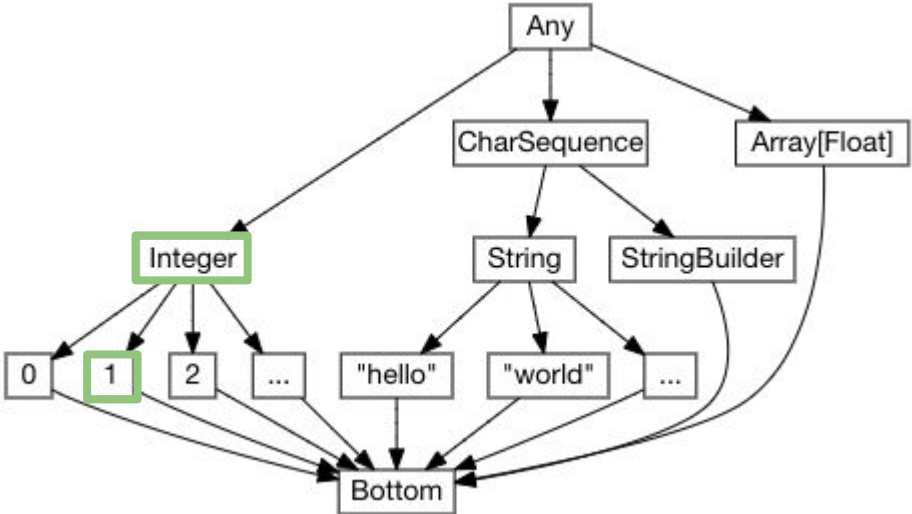


Recursive Inter-Procedural Inference

```

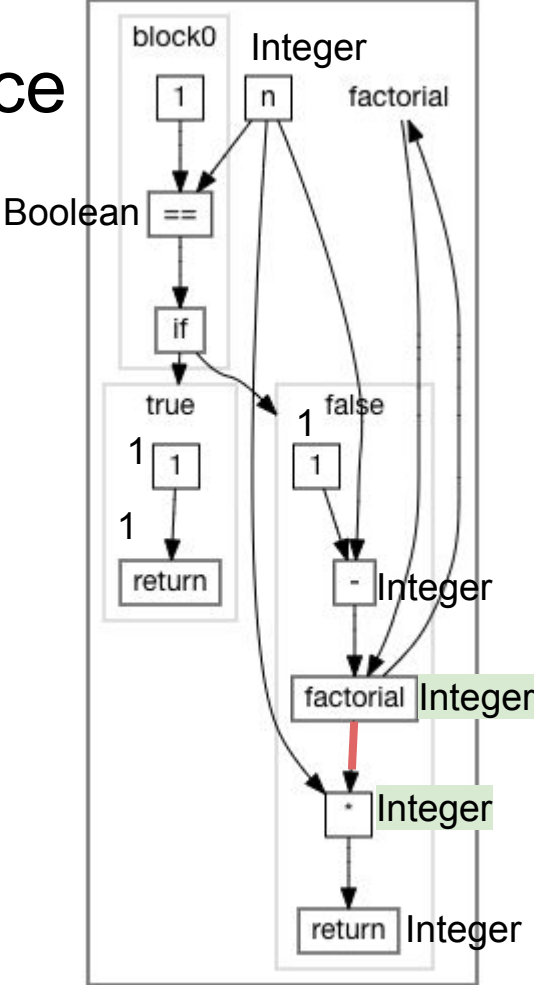
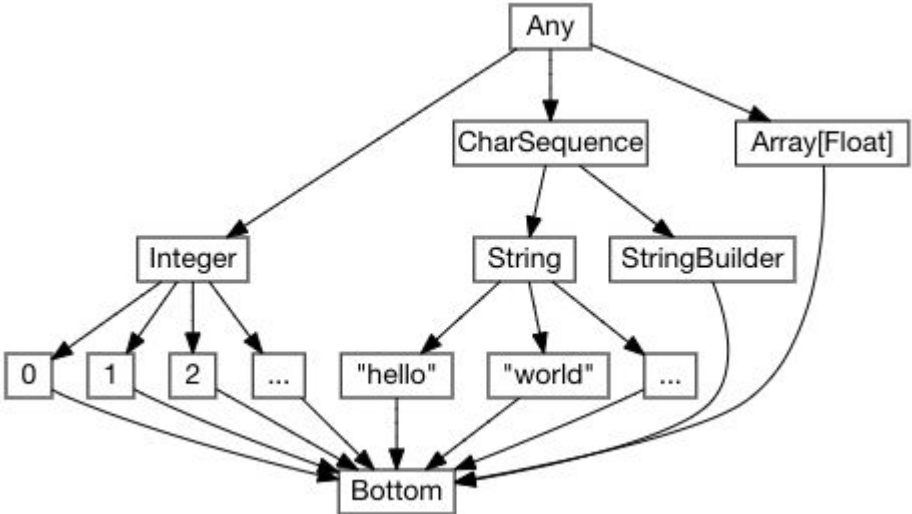
public static Any factorial(int n) {
    if (n == 1) {
        return 1;
    } else {
        return n * factorial(n - 1);
    }
}

```



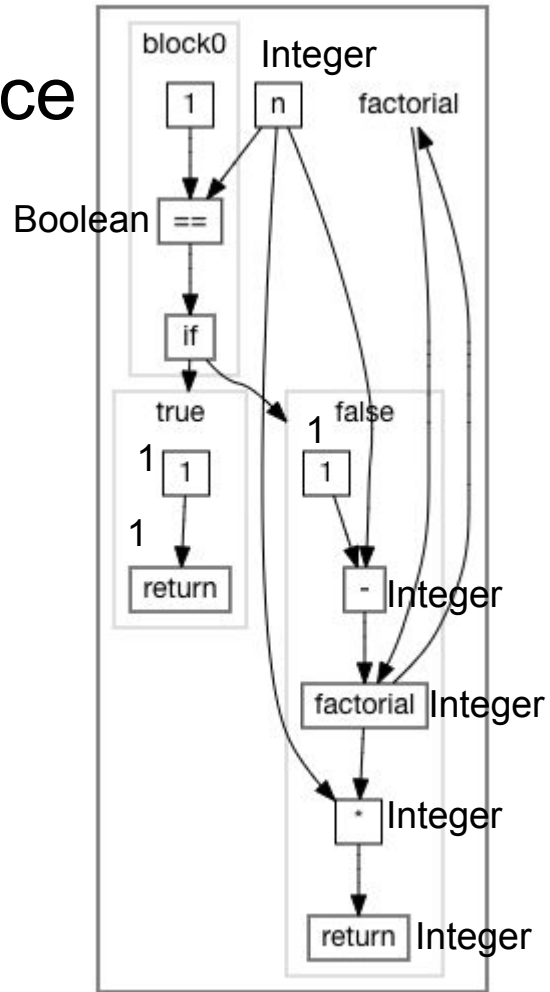
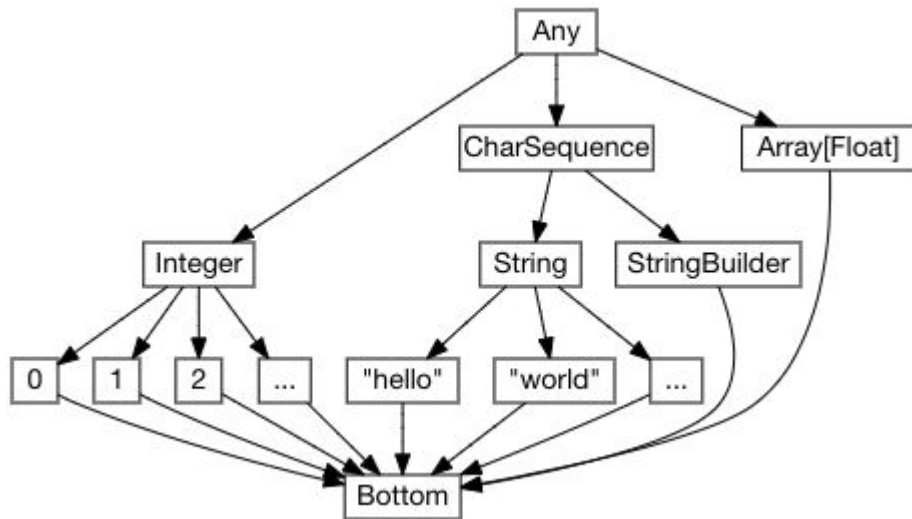
Recursive Inter-Procedural Inference

```
public static Any factorial(int n) {  
    if (n == 1) {  
        return 1;  
    } else {  
        return n * factorial(n - 1);  
    }  
}
```



Recursive Inter-Procedural Inference

```
public static Any factorial(int n) {  
public static Integer factorial(int n) {  
    if (n == 1) {  
        return 1;  
    } else {  
        return n * factorial(n - 1);  
    }  
}
```



How an Optimizing Compiler Works

Hand Optimizing Some Code

Modelling a Program

Inferences and Optimizations

- Type Inference & Constant Folding
- Inter-Procedural Inference
- Recursive Inter-Procedural Inference
- **Liveness & Reachability Analysis**

Liveness & Reachability Analysis

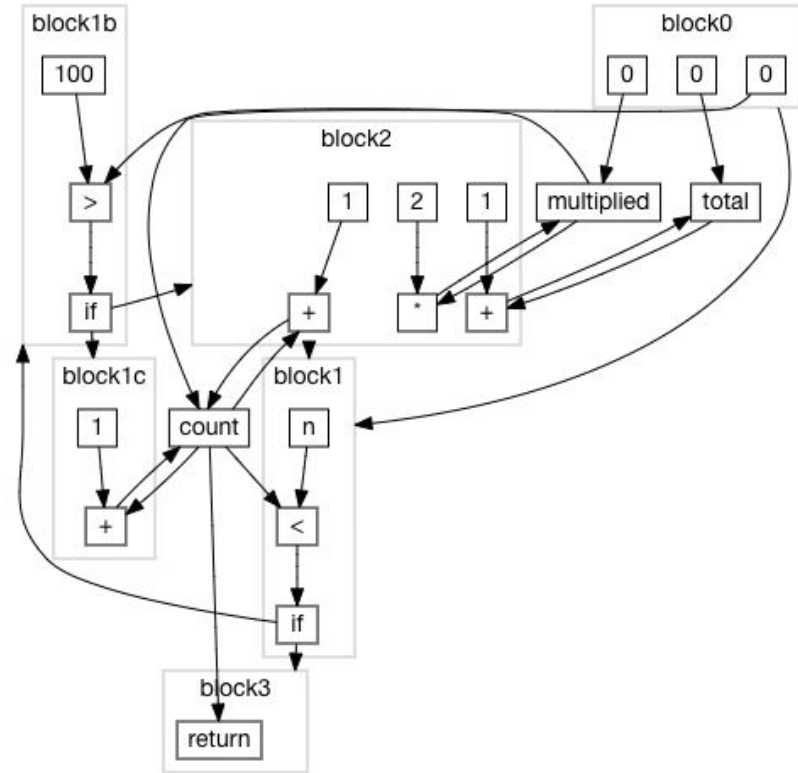
Find all the code whose values contribute to the final returned result ("Live")

Find all code which the the control-flow of the program can reach ("Reachable")

Code that fails either test is a safe candidate for removal!

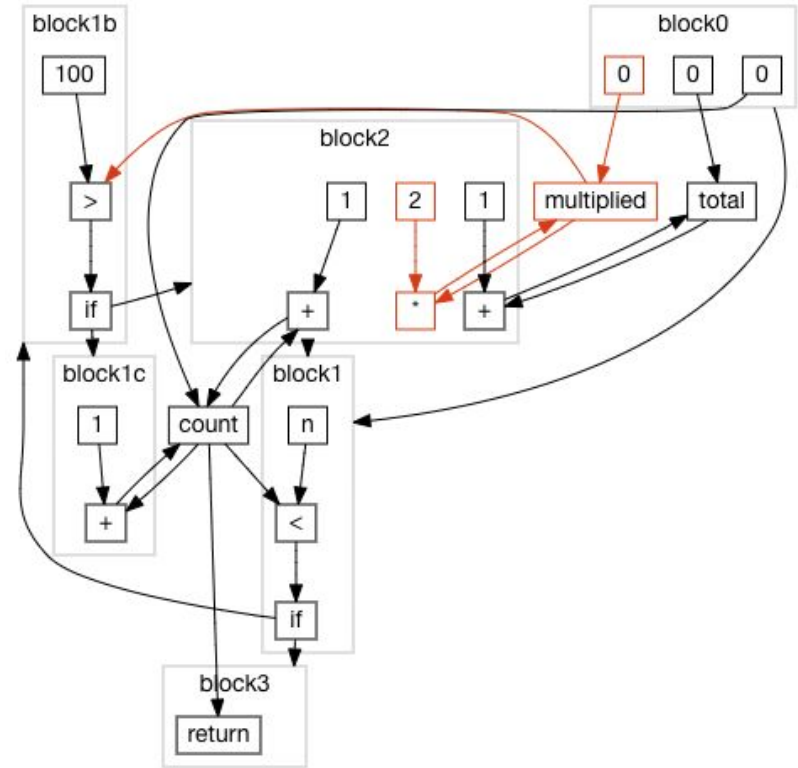
Strawman Program

```
static int main(int n){  
    int count = 0, total = 0, multiplied = 0;  
    while(count < n){  
        if (multiplied > 100) count += 1;  
        count += 1;  
        multiplied *= 2;  
        total += 1;  
    }  
    return count;  
}
```



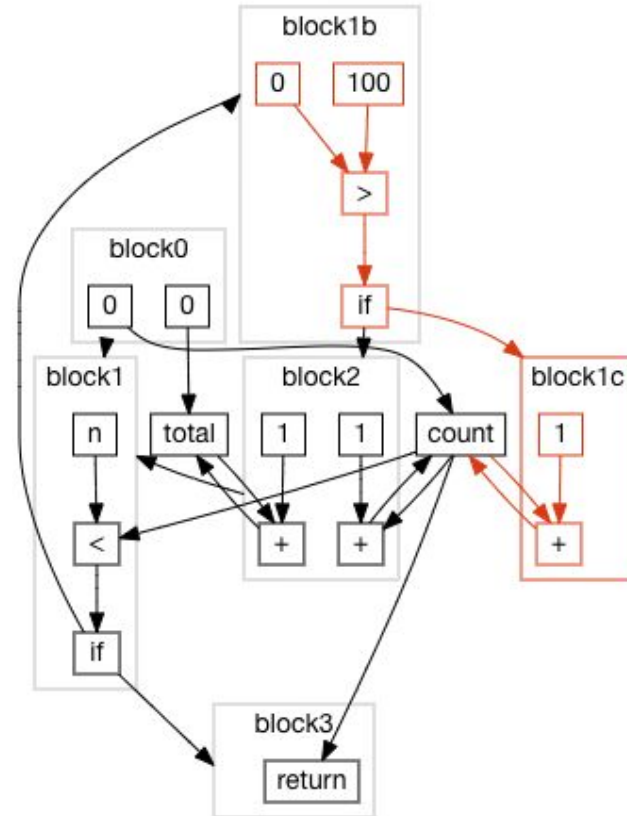
Type Inference & Constant Folding

```
static int main(int n){
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  while(count < n){
    if (multiplied > 100) count += 1;
    count += 1;
    multiplied *= 2;
    total += 1;
  }
  return count;
}
```



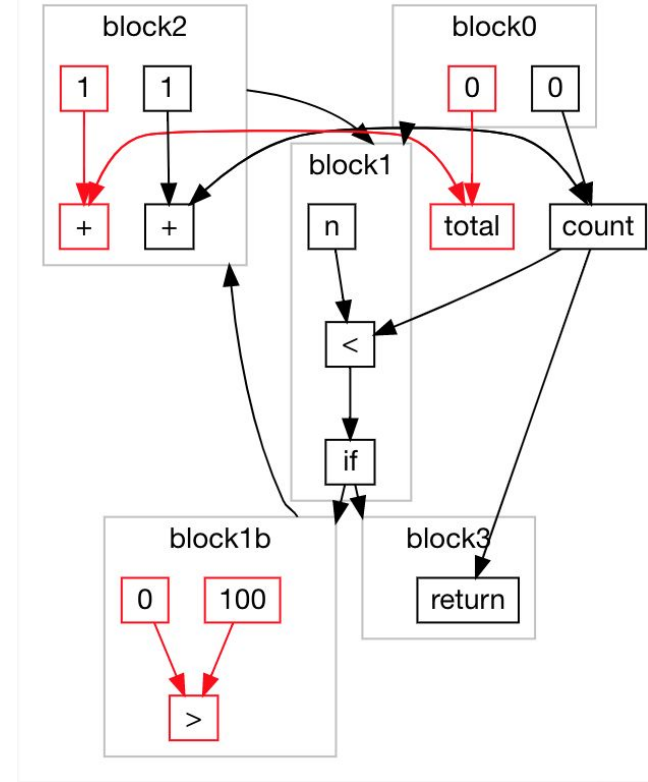
Branch Elimination & Reachability Analysis

```
static int main(int n){  
    int count = 0, total = 0;  
    while(count < n){  
        if (0 > 100) count += 1;  
        count += 1;  
        total += 1;  
    }  
    return count;  
}
```



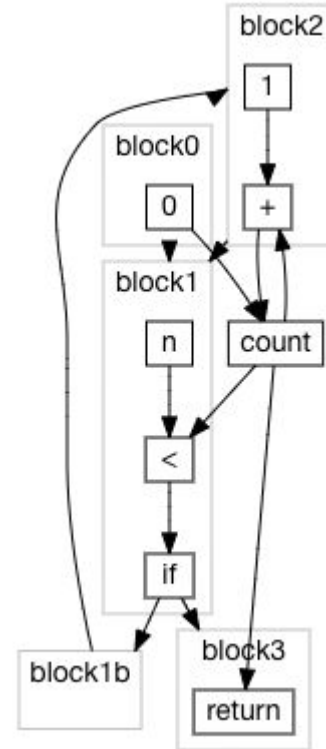
Liveness Optimizations

```
static int main(int n){  
    int count = 0, total = 0;  
    while(count < n){  
        0 > 100;  
        count += 1;  
        total += 1;  
    }  
    return count;  
}
```



Final Output Code

```
static int main(int n){  
    int count = 0;  
    while(count < n){  
        count += 1;  
    }  
    return count;  
}
```



How an Optimizing Compiler Works

Hand Optimizing Some Code

- Type Inference
- Inlining
- Constant Folding
- Dead Code Elimination
- Branch Elimination
- Late Scheduling

Modelling a Program

- Sourcecode
- Abstract Syntax Trees
- Bytecode
- Dataflow Graphs

Making Inferences and Optimizations

- Type Inference & Constant Folding
- Inter-Procedural Inference
- Recursive Inter-Procedural Inference
- Liveness & Reachability Analysis

How an Optimizing Compiler Works

Hand Optimizing Some Code

- Type Inference
- Inlining
- Constant Folding
- Dead Code Elimination
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- Late Scheduling

[Engineering a Compiler by Keith D Cooper & Linda Torczon](#)

[Combining Analyses, Combining Optimizations by Cliff Click](#)

Modelling a Program

- Sourcecode
- Abstract Syntax Trees
- Bytecode
- Dataflow Graphs

Making Inferences and Optimizations

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